

**INTRODUCTION TO  
COMPUTATIONAL  
TOPOLOGY**

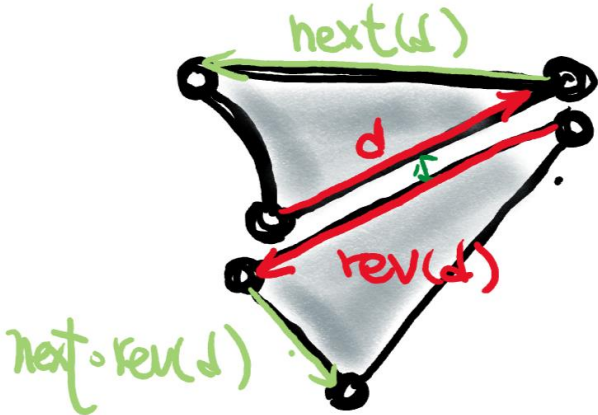
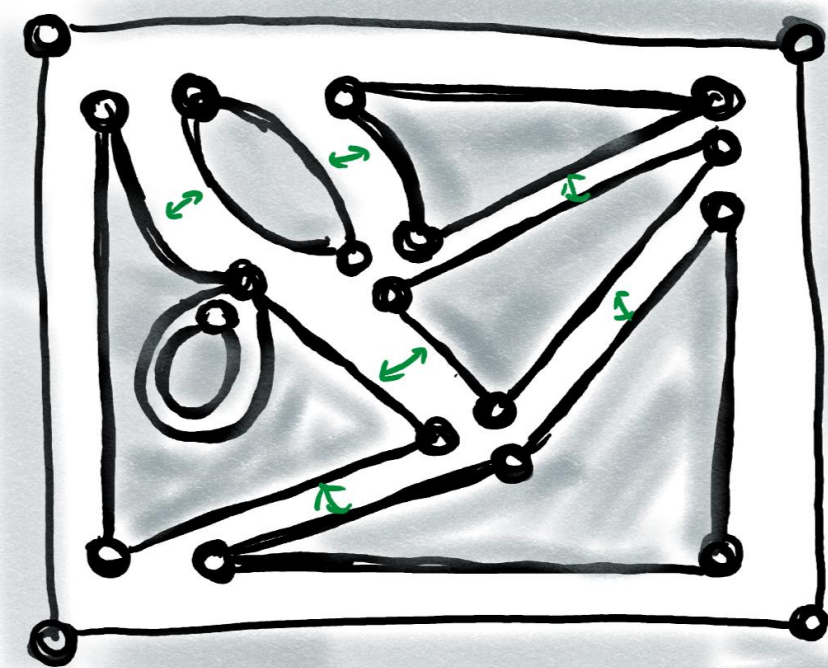
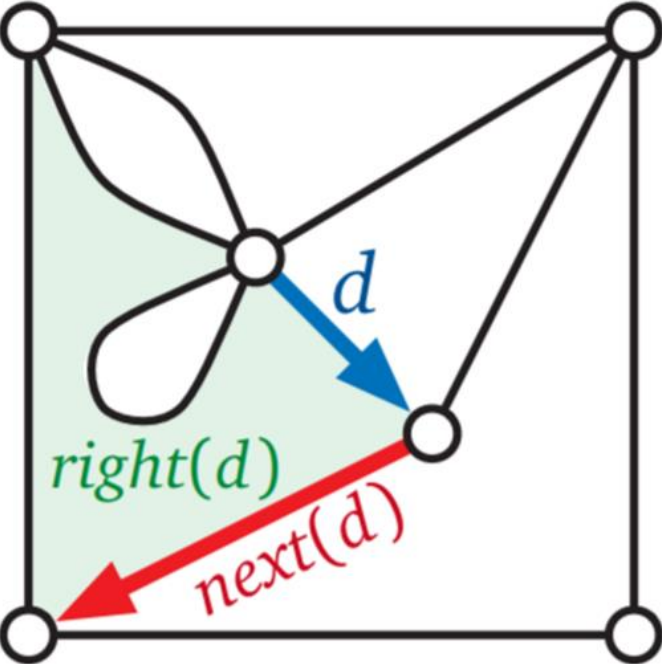
**HSIEN-CHIH CHANG  
LECTURE 4, SEPTEMBER 23, 2021**

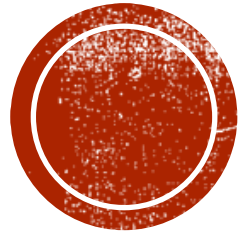
# ACKNOWLEDGEMENT

- Most of the figures today comes from
  - Jeff Erickson, *One-Dimensional Computational Topology*
  - Robert Christ, *Elementary Applied Topology*
  - Keenan Crane, *Discrete differential geometry: An applied introduction*



# RECAP: POLYGONAL SCHEMA IS ROTATION SYSTEM

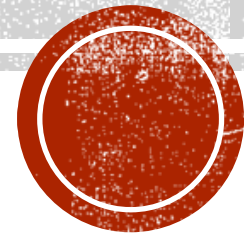




# EULER'S FORMULA

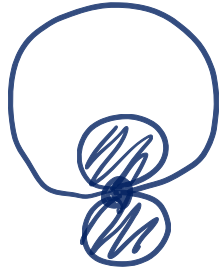


**LET'S FOCUS ON PLANE GRAPHS**



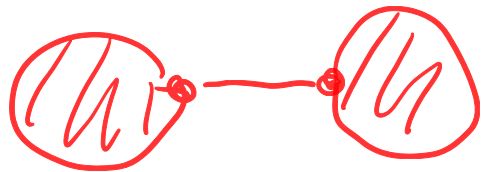
# SPANNING TREE TERMINOLOGY

## ■ Loop



edge whose two incident vertices are the same

## ■ Bridge

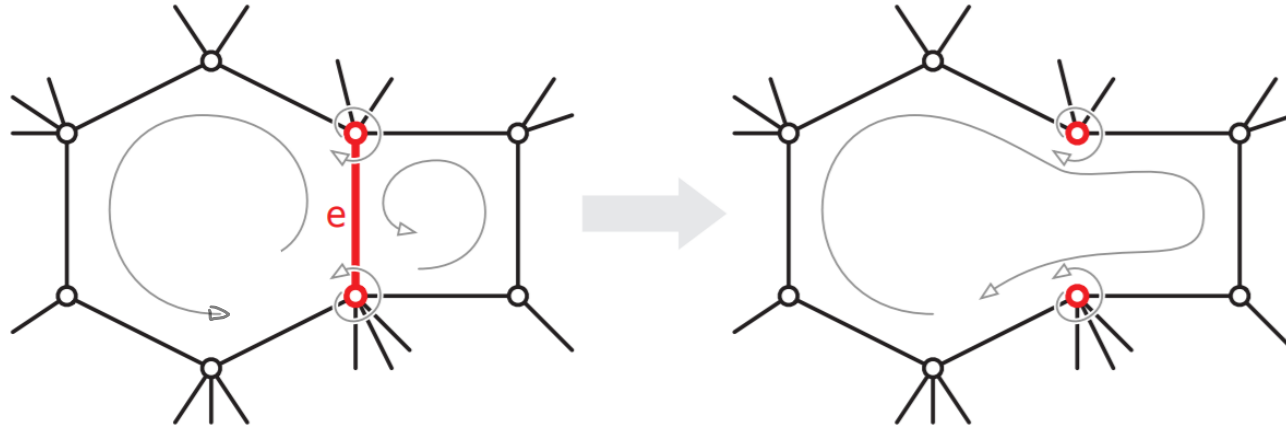


edge whose two incident faces are the same

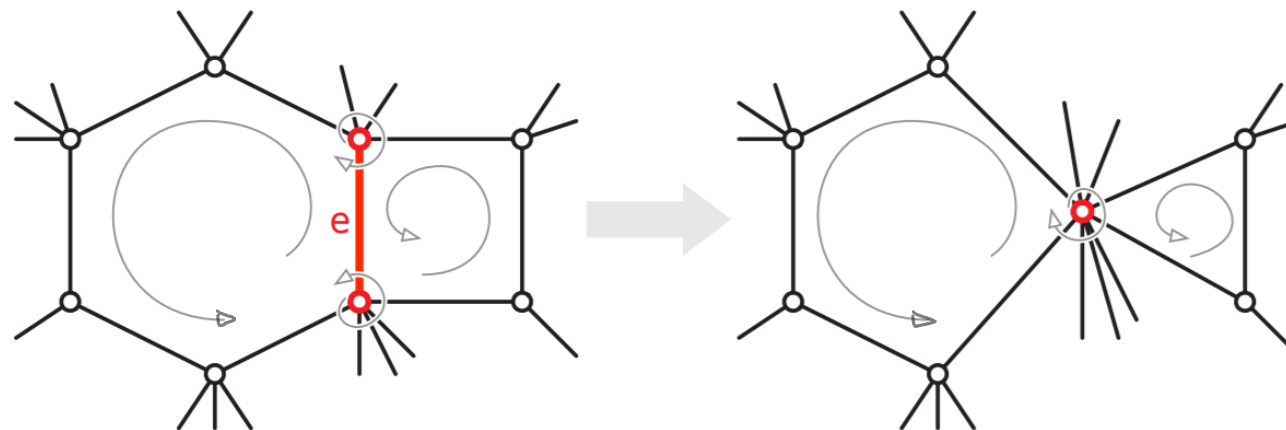


# SPANNING TREE TERMINOLOGY

■ Delete



■ Contract



# THE SPANNING TREE ALGORITHM

SPANNING TREE (G):

for any edge  $e \in G$ :

if  $e$  is a loop:

delete  $e$

if  $e$  is a bridge:

contract  $e$

otherwise:

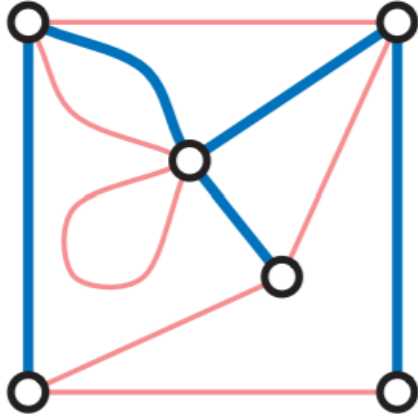
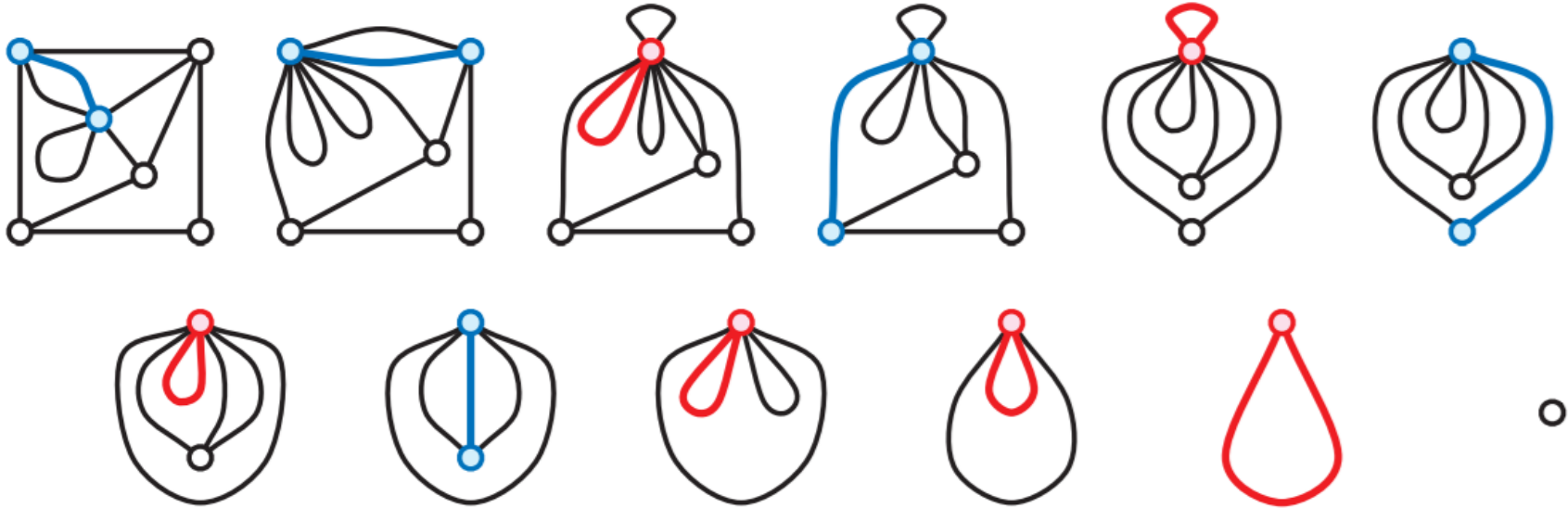
delete or contract  $e$

return all contracted edge  $T$





# THE SPANNING TREE ALGORITHM



Computing a spanning tree of a graph.



# THE SPANNING TREE ALGORITHM

WHATEVERFIRSTSEARCH( $s$ ):

put  $(\emptyset, s)$  in bag

while the bag is not empty

    take  $(p, v)$  from the bag (★)

    if  $v$  is unmarked

        mark  $v$

$parent(v) \leftarrow p$

        for each edge  $vw$  (†)

            put  $(v, w)$  into the bag (★★)



# WHAT HAPPENED IN THE DUAL?

SPANNING TREE ( $G$ ):

for any edge  $e$  in  $G$  :  
if  $e$  is a loop:  
delete  $e$ .  
if  $e$  is a bridge:  
contract  $e$   
o.w.  
delete or contract  $e$   
return all contracted edges  $T$

SPANNING TREE ( $G^*$ ):

for any edge  $e^*$  in  $G^*$  :  
if  $e^*$  is a ~~loop~~ <sup>bridge</sup> :  
~~delete~~ <sup>contract</sup>  $e^*$   
if  $e^*$  is a ~~bridge~~ <sup>loop</sup> :  
~~contract~~ <sup>delete</sup>  $e^*$   
o.w. ~~delete~~ <sup>contract</sup> or ~~contract~~ <sup>delete</sup>  $e^*$   
return all contracted edges  $T$

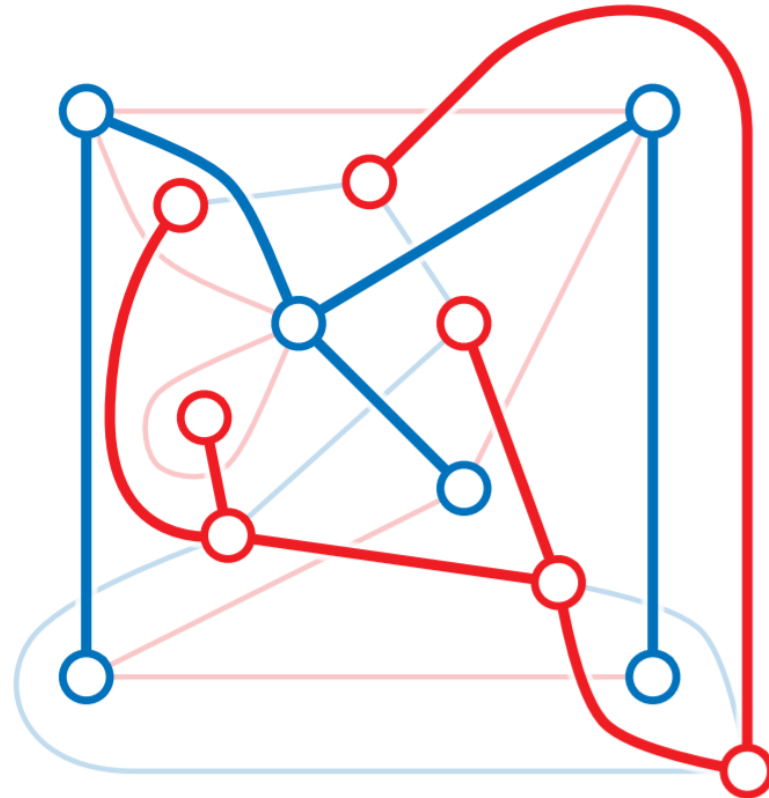


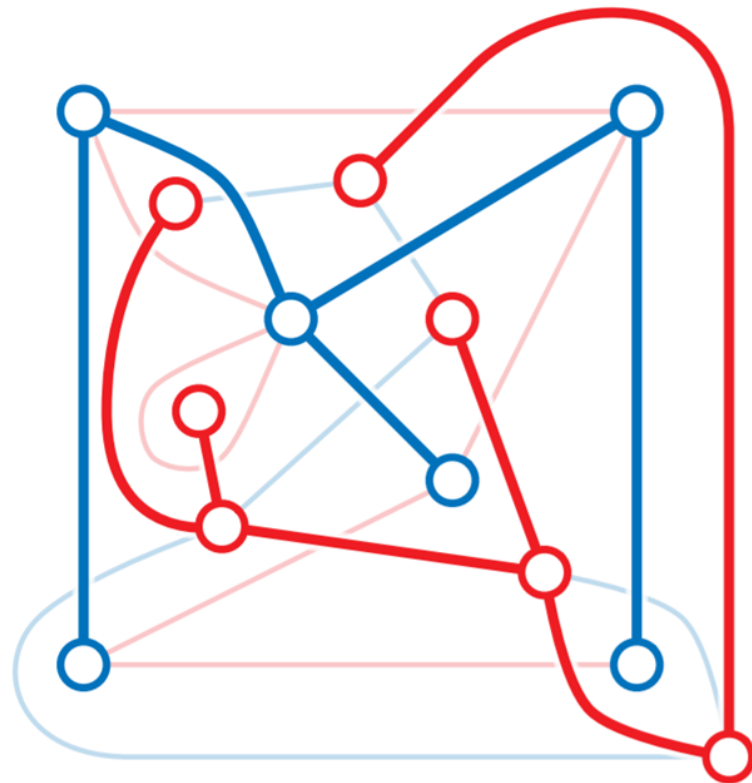
# TREE-COTREE DECOMPOSITION

- Plane graph  $G$  decomposes into
  - Primal spanning tree  $T$
  - Dual spanning cotree  $C$

$$E = (V - 1) + (F - 1)$$

$$V - E + F = 2$$



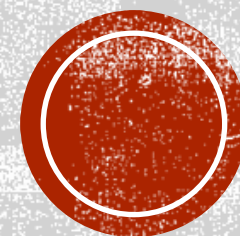


# EULER'S FORMULA

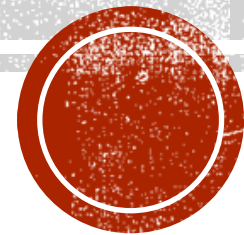
[Euler 1750] [Legendre 1794] [Cayley-Listing 1861]

For any plane graph  $G$ ,

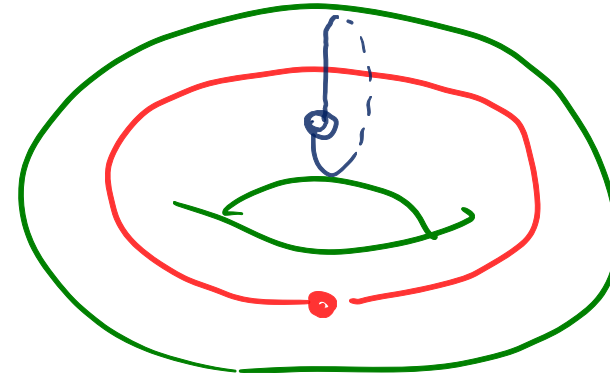
$$V_G - E_G + F_G = 2$$



# WHAT ABOUT SURFACE GRAPH?



# SPANNING TREE



SPANNING TREE ( $G$ ):

for any edge  $e$  in  $G$ :

if  $e$  is a loop w/ diff. incident faces  
delete  $e$ .

if  $e$  is a bridge: has diff. endpoints,  
contract  $e$

o.w.

delete or contract  $e$

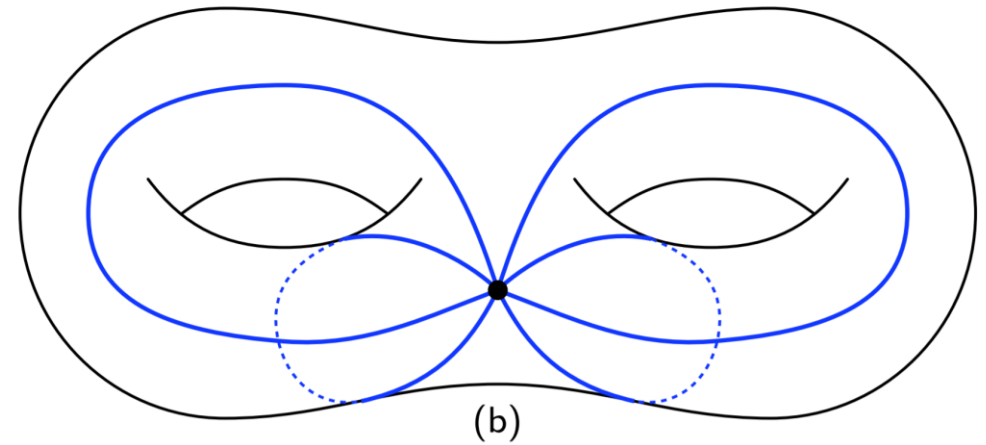
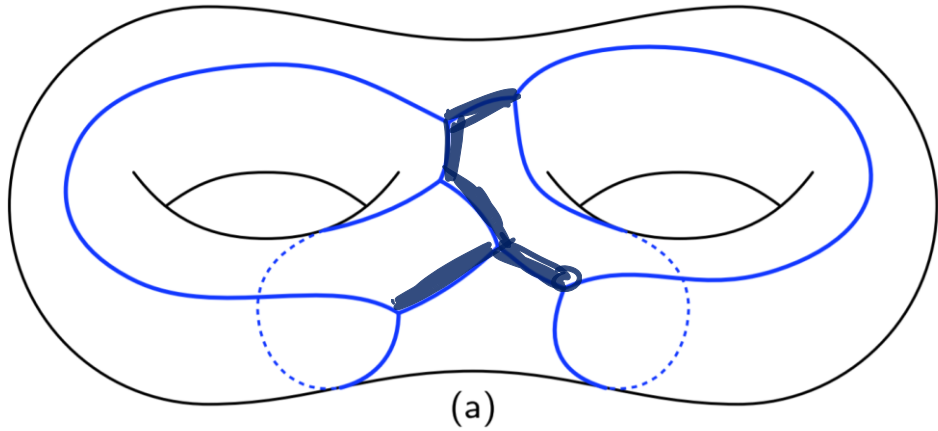
return all contracted edges  $T$



# SYSTEM OF LOOPS

Euler  
Characteristic

$$E = (V-1) + (F-1) + L$$
$$\chi = V - E + F = 2 - L$$

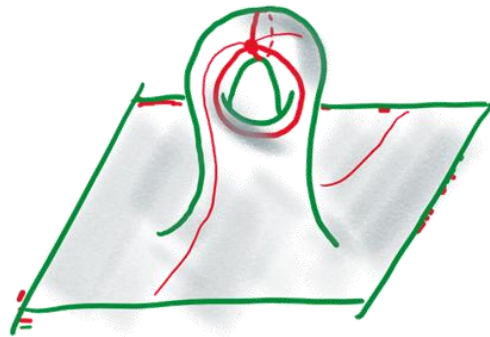


leftovers  $L$

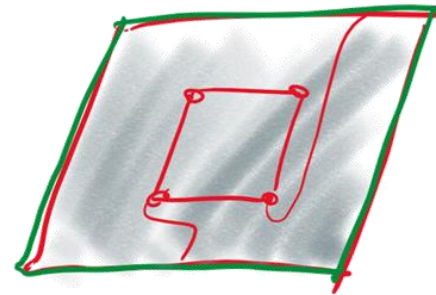
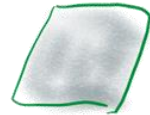
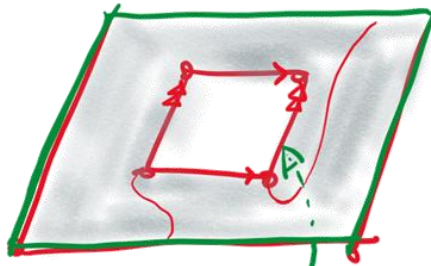




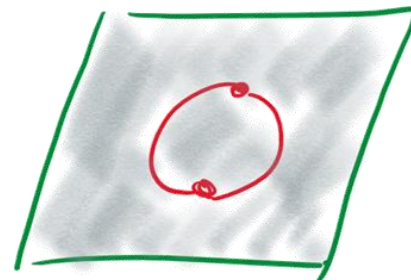
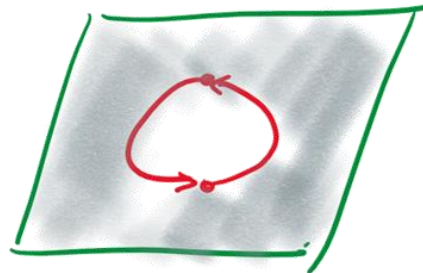
# HOW MANY LEFTOVER EDGES?



$(V, E, F) \chi$



$(V+3, E+2, F+1) \chi+2$

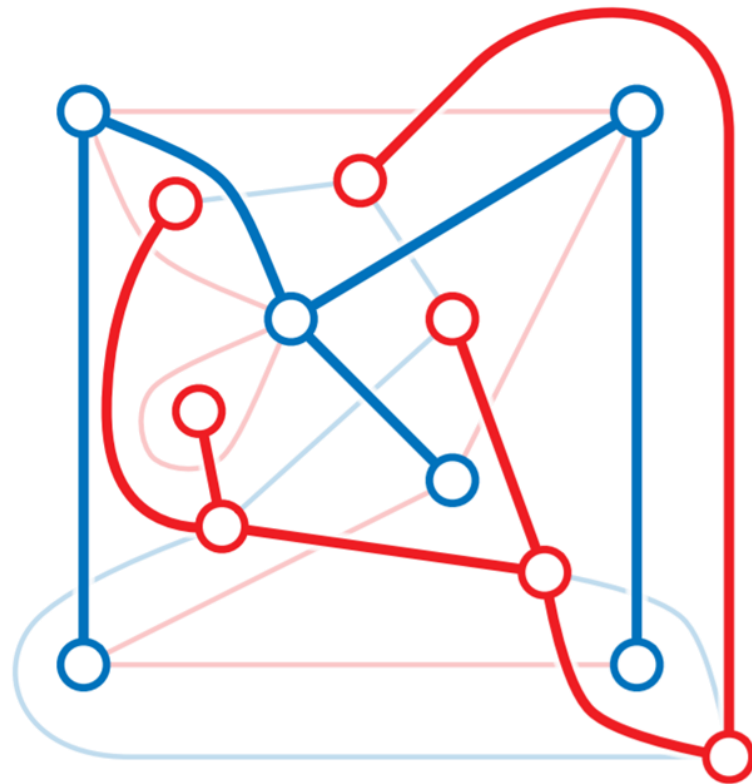


$(V+1, E+1, F+1) \chi+1$

$g$  handle  
 $\Rightarrow \chi + 2g$

$r$  crosscaps  
 $\Rightarrow \chi + r$



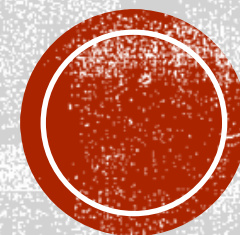


# EULER'S FORMULA

[Euler 1750] [Legendre 1794] [Cayley-Listing 1861]

For any graph  $G$  embedded on surface  $\Sigma(g,r,b)$ ,

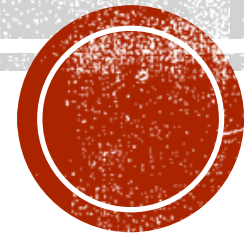
$$V_G - E_G + F_G = 2 - 2g - r - b$$



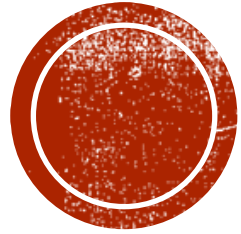
# EULER CHARACTERISTIC IS A “COMPLETE” INVARIANT OF SURFACES

**PONDER.**

Torus and Möbius band  
have the same  $\chi$ ?



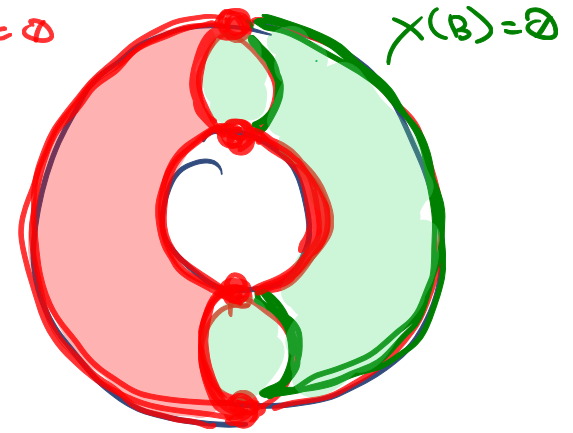
# INTERMISSION



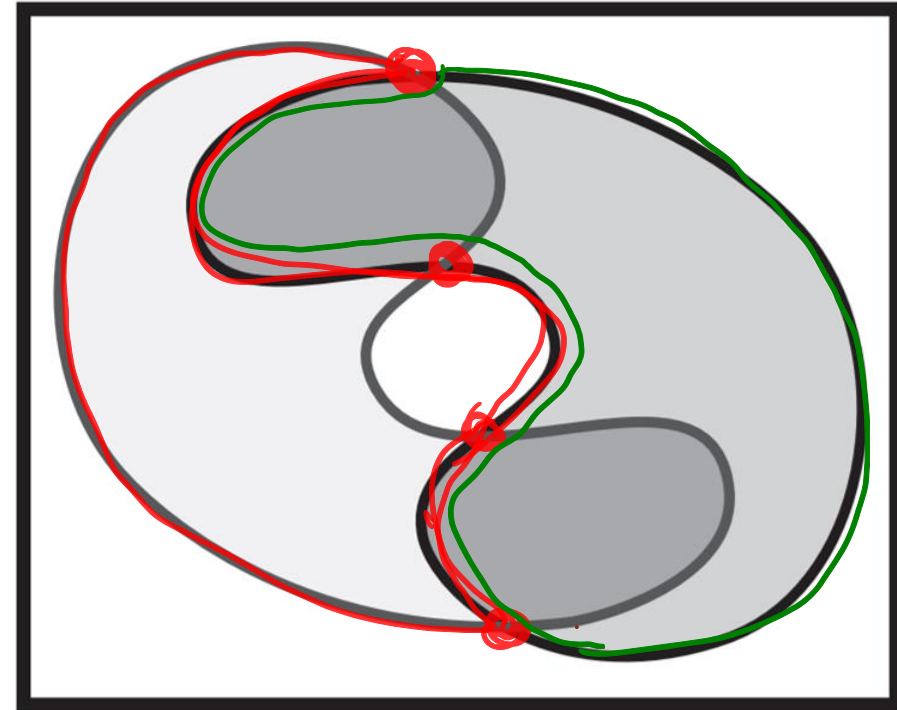
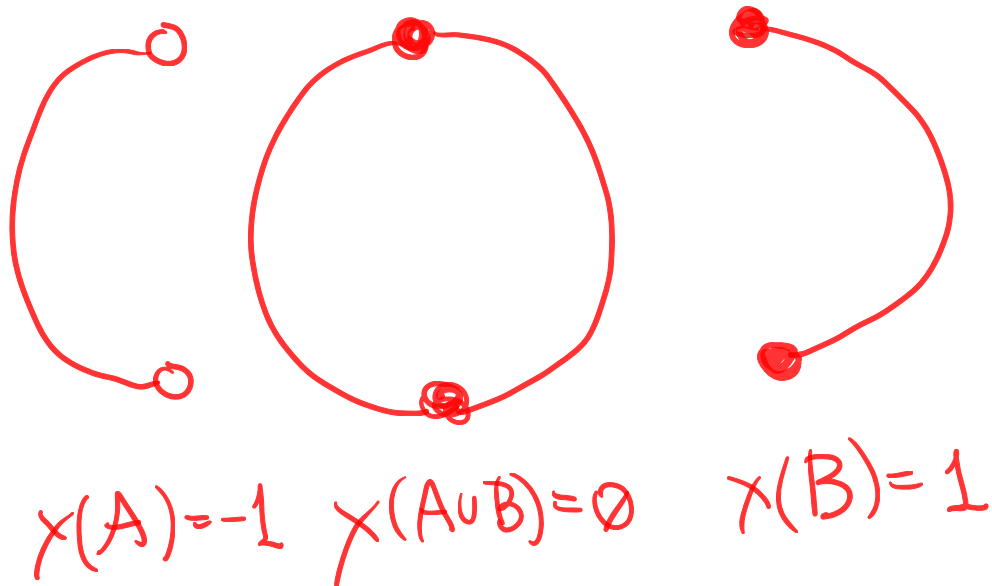
# EULER CALCULUS AND CURVATURE

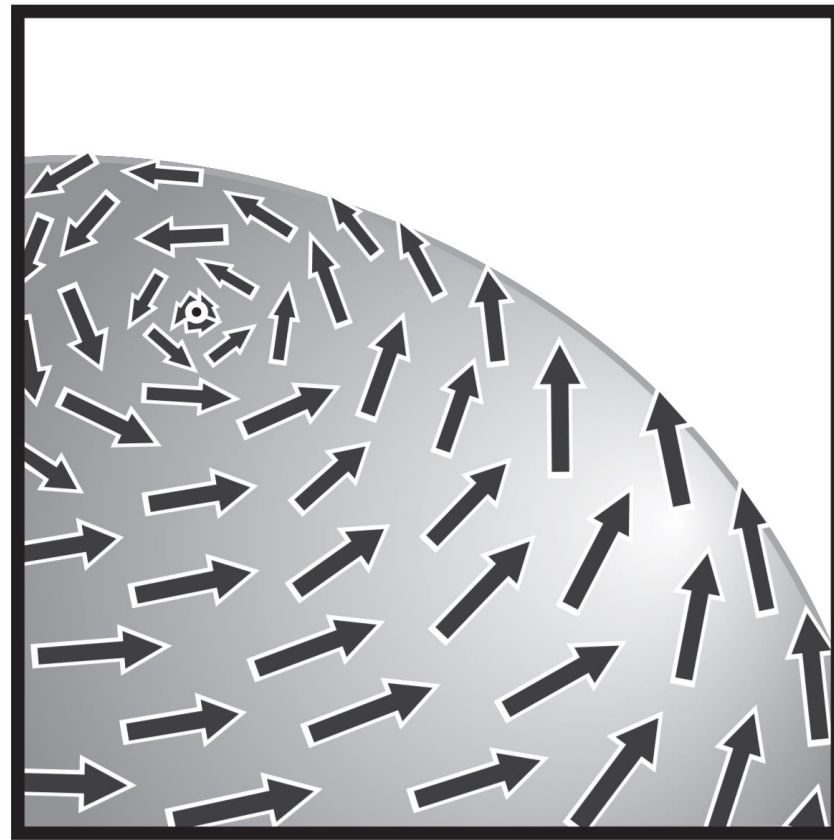


# EULER CHARACTERISTIC IS ADDITIVE!



- $\chi(A \cup B) = \chi(A) + \chi(B) - \chi(A \cap B)$
- Watch out for open/closed

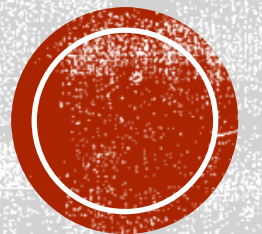




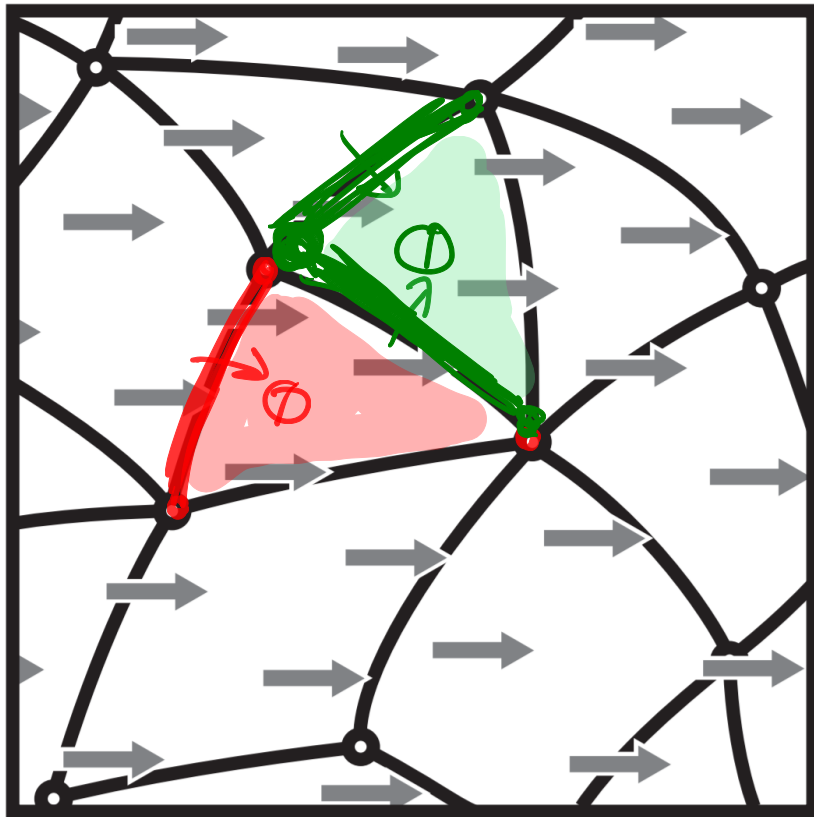
## HEDGEHOG THEOREM

[Poincaré 1885] [Brouwer 1912]

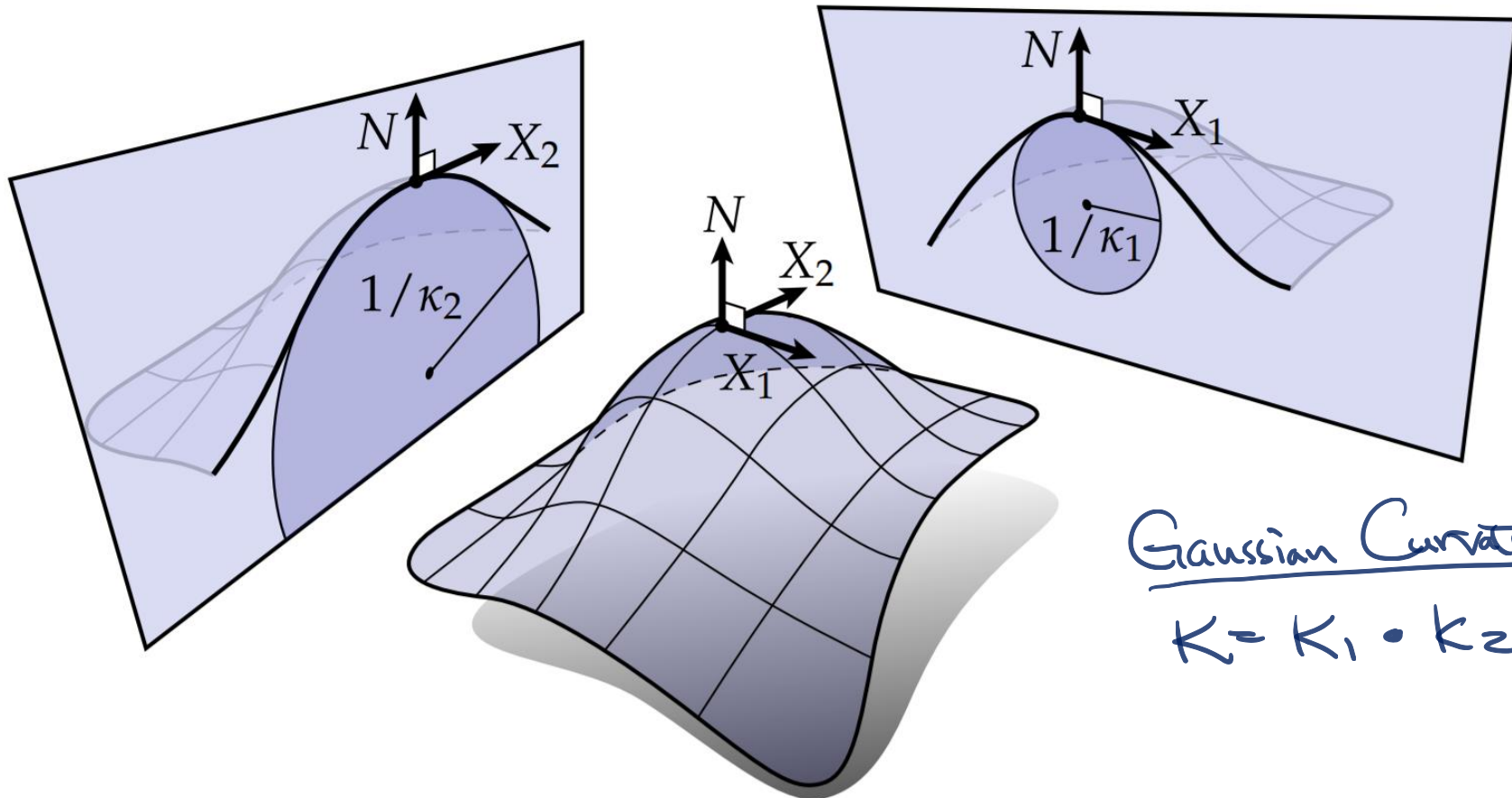
There is no non-vanishing continuous vector field on closed surfaces of non-zero Euler characteristics.



# EULER CALCULUS PROOF



# CURVATURE



Gaussian Curvature

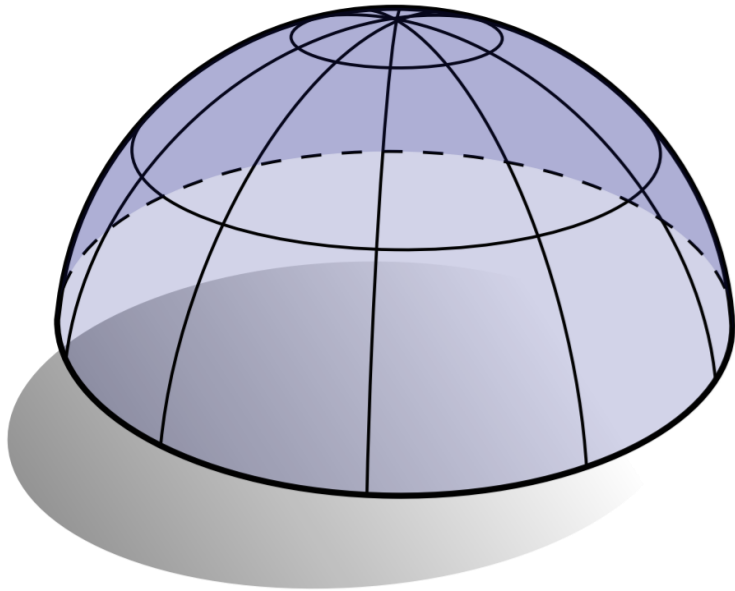
$$K = \kappa_1 \cdot \kappa_2$$



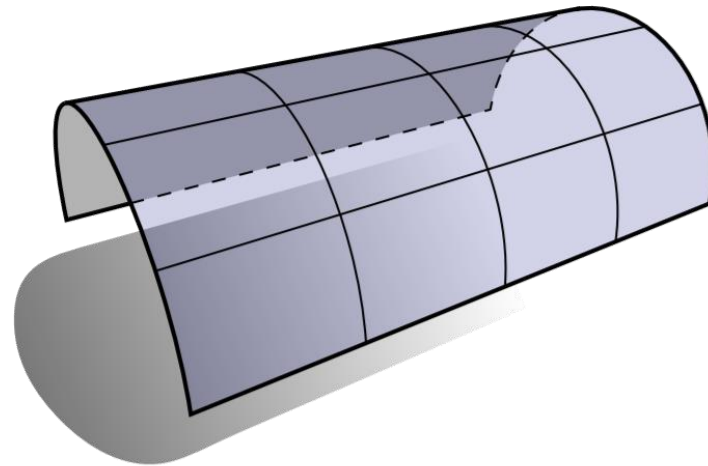


# CURVATURE

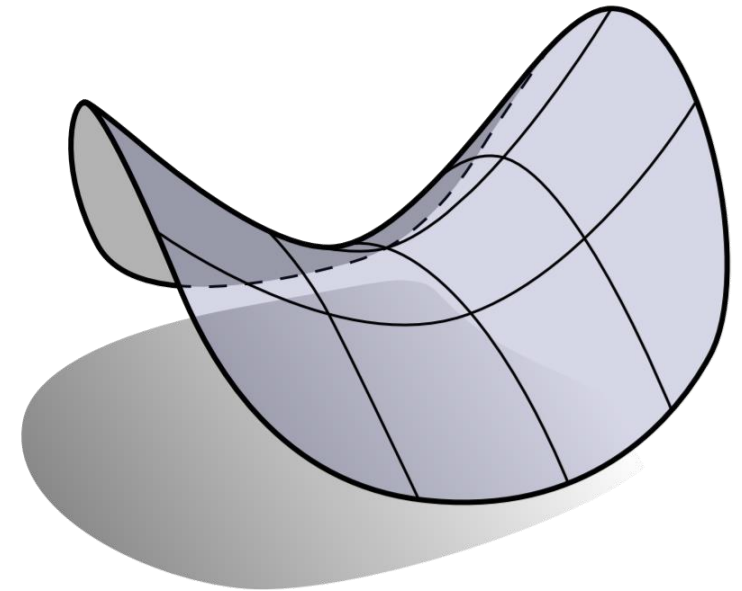
$$K > 0$$

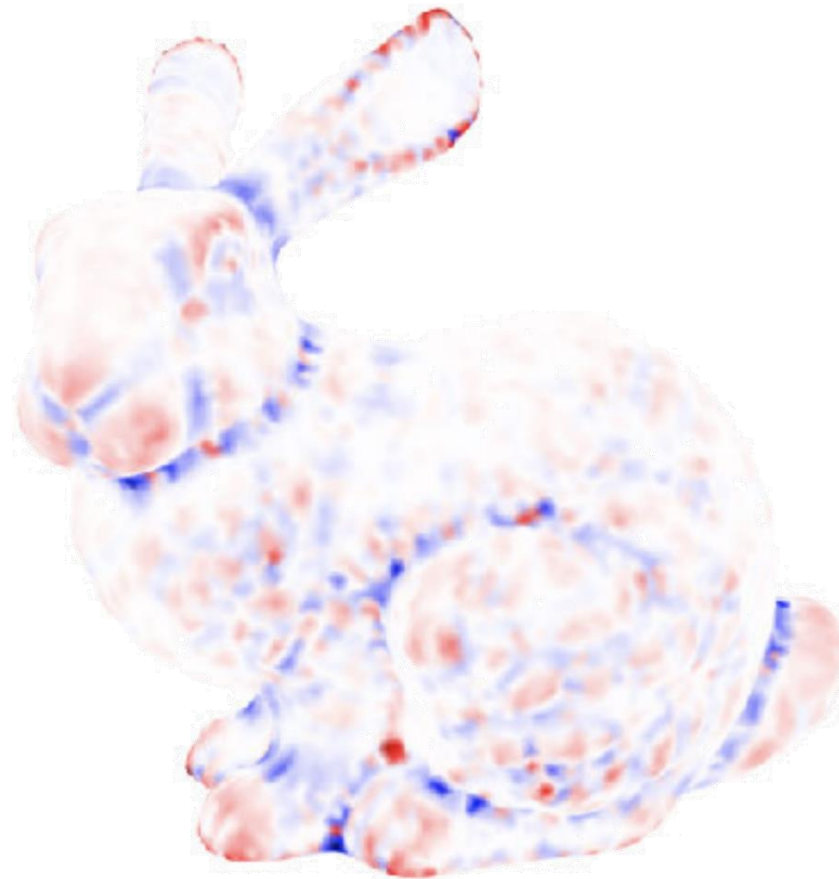
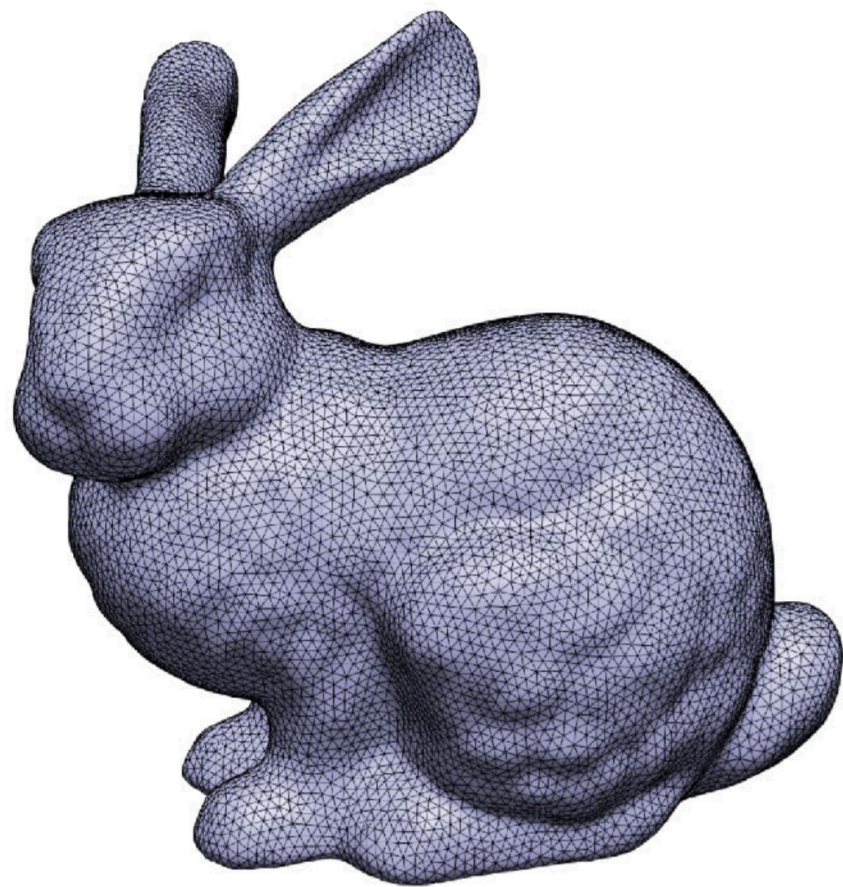


$$K = 0$$



$$K < 0$$

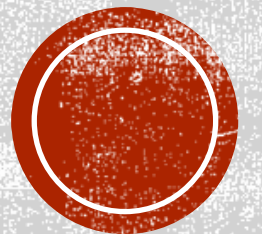




# GAUSS-BONNET THEOREM [Gauß 1827] [Bonnet 1848]

For any surface  $\Sigma$  of Euler char.  $\chi$  and curvature  $\kappa$ ,

$$\int_{\Sigma} \kappa \, d\sigma = 2\pi\chi$$

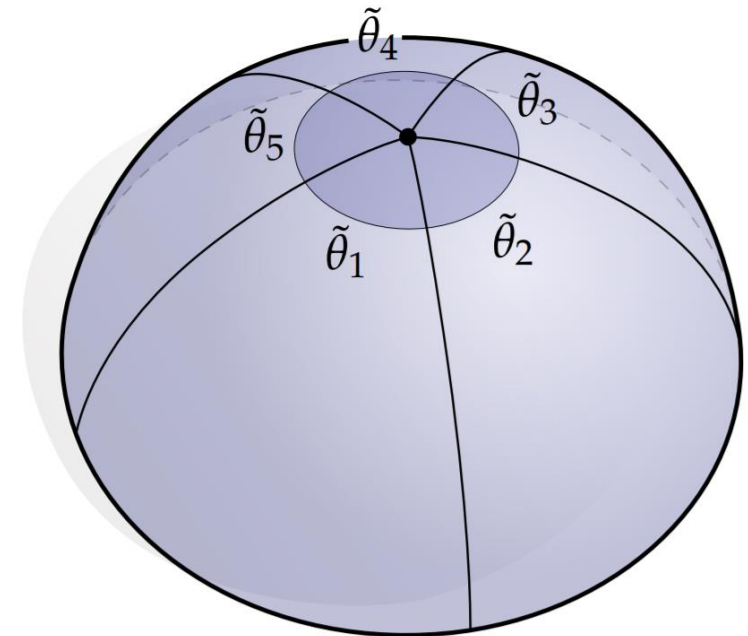
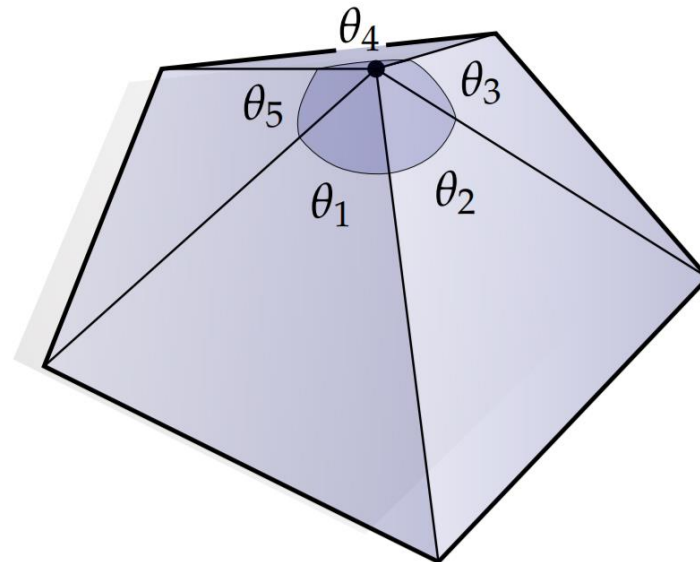


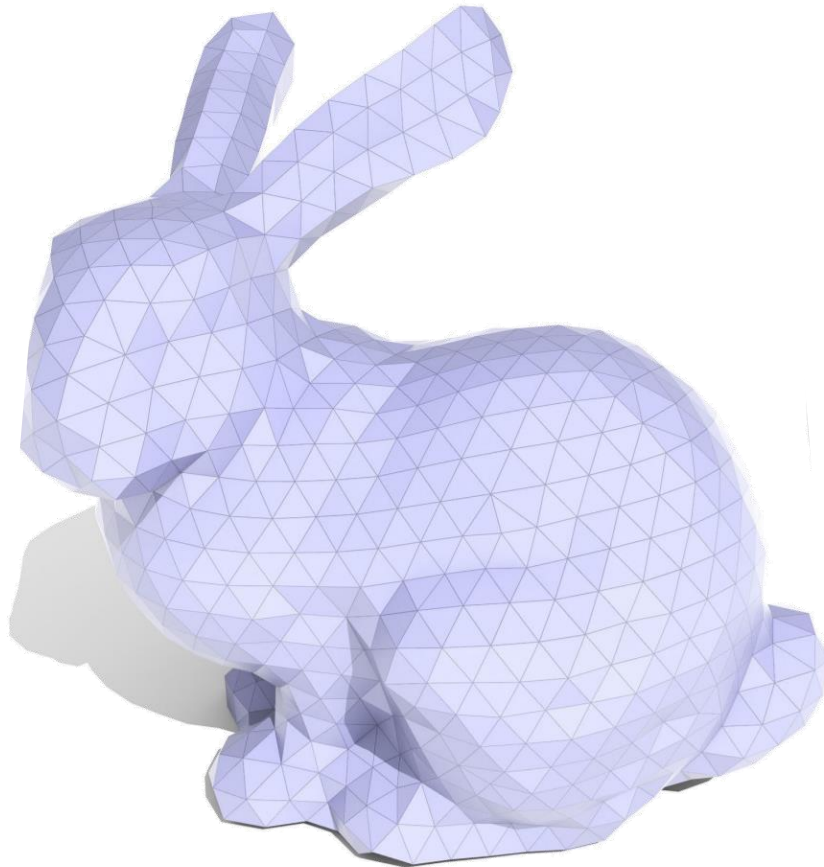
# CURVATURE

- Curvature  $\kappa_\theta(x)$

$$\kappa_\theta(x) = 2\pi - \sum_i \theta_i$$

angle defect

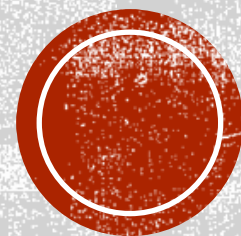




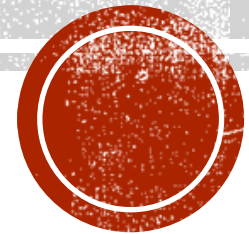
## DISCRETE GAUSS-BONNET THEOREM

For any discrete surface  $\Sigma$  of Euler char.  $\chi$  and angles  $\theta$ ,

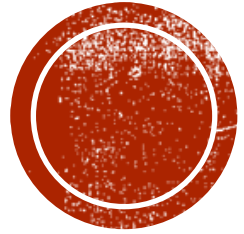
$$\sum_x \kappa_\theta(x) = 2\pi\chi$$



# **CURVATURES CAN BE MOVED AROUND, BUT NOT REMOVED**

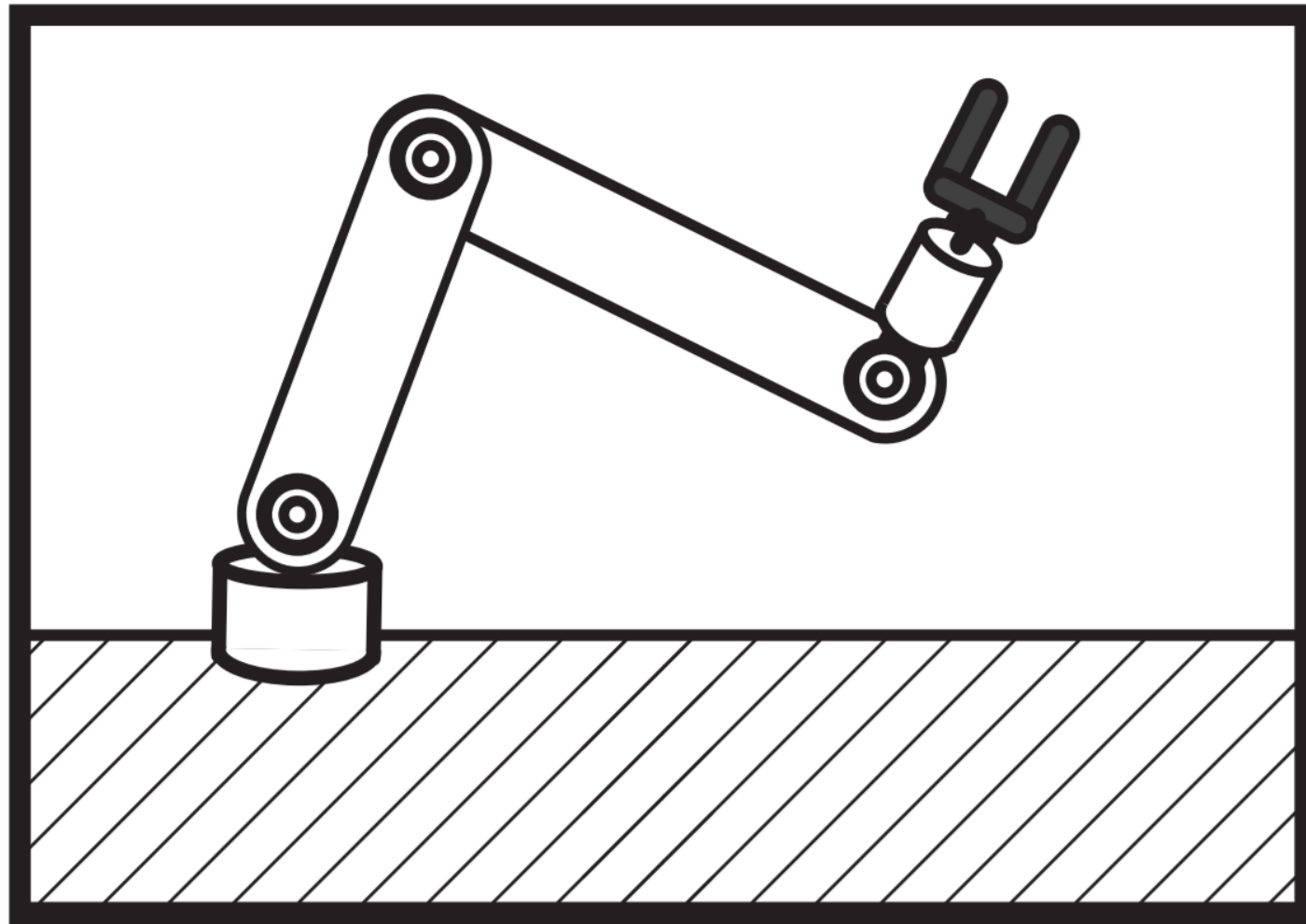


**TO THINK ABOUT LATER.  
Can you prove Hedgehog  
Theorem using Gauss-Bonnet?**



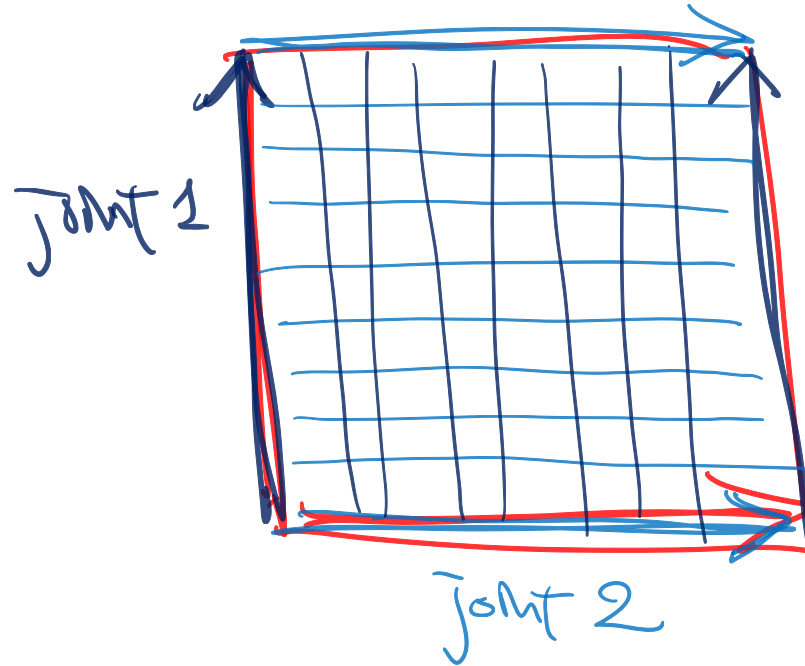
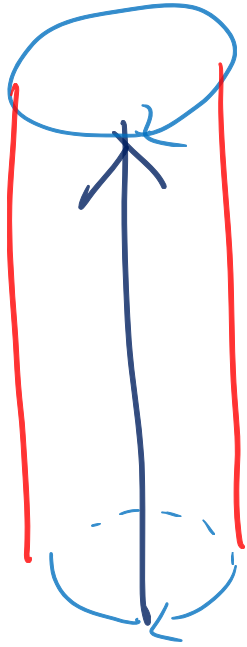
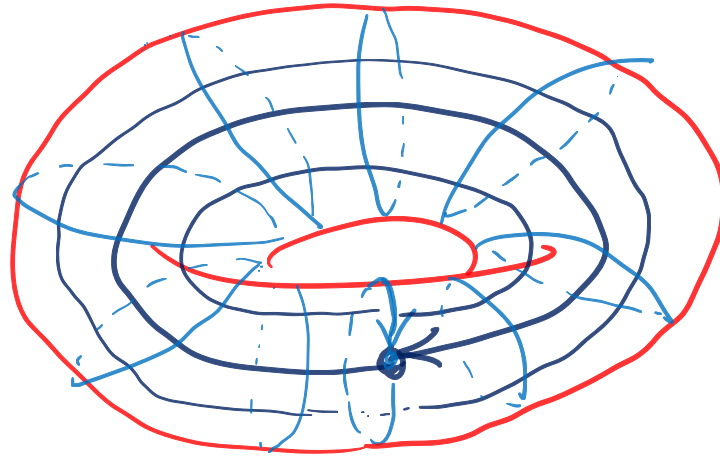
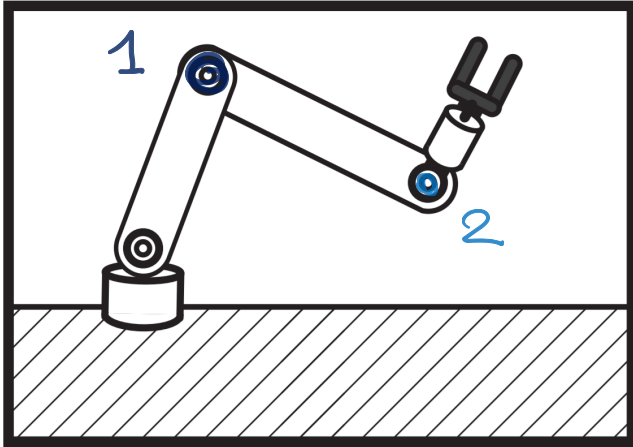
# CONFIGURATION SPACE AND COMPLEX





# EXAMPLE: ROBOT ARMS

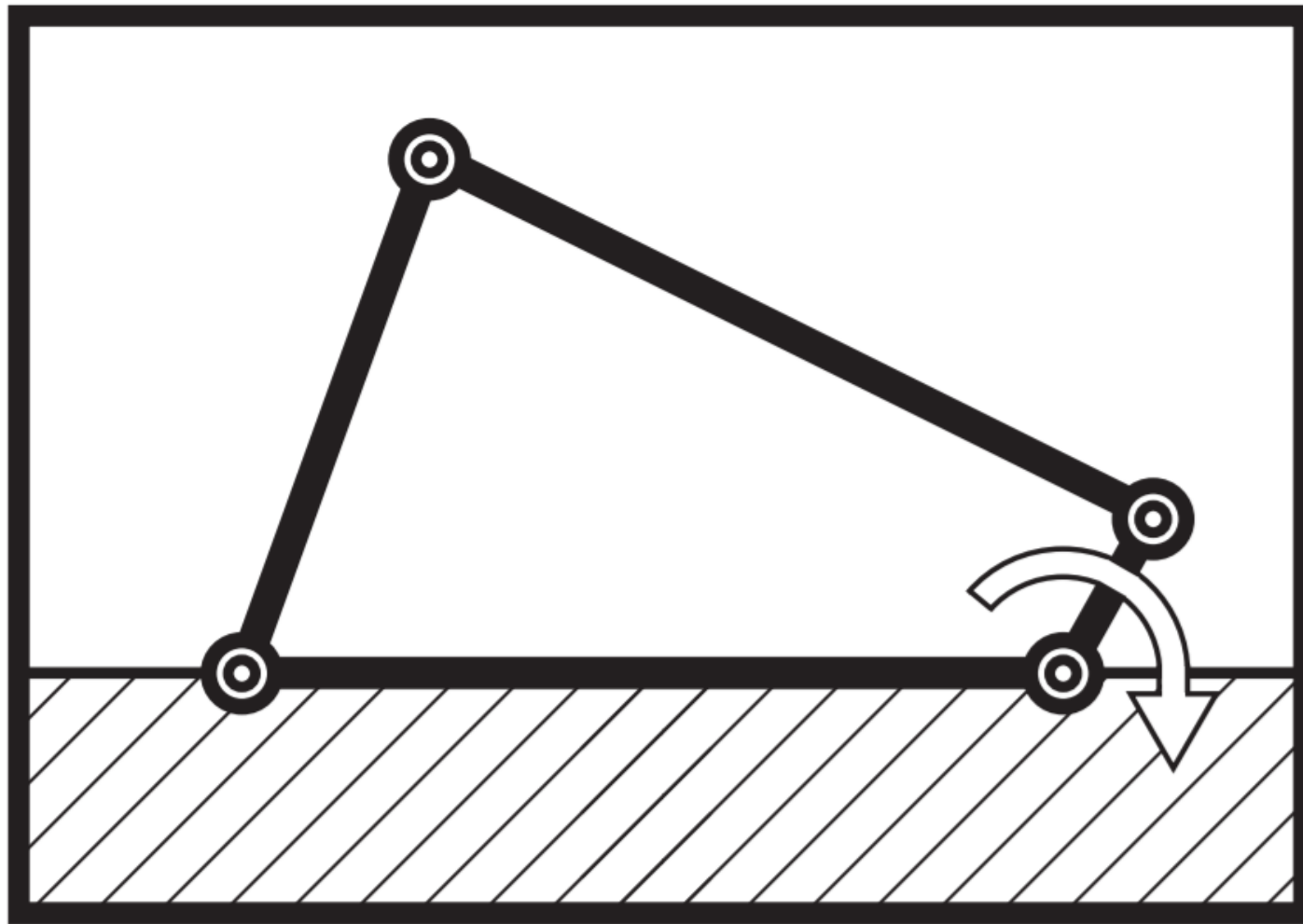




# EXAMPLE: ROBOT ARMS

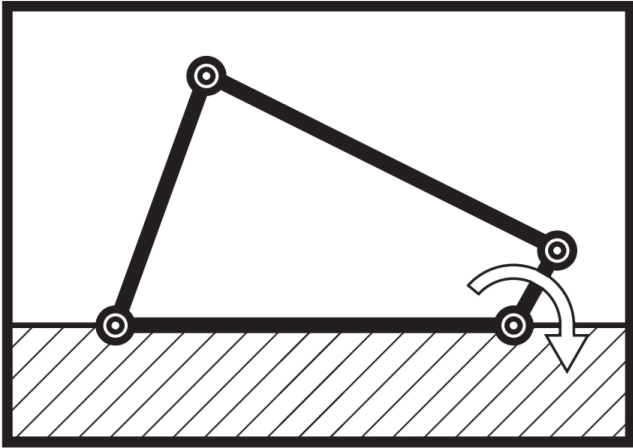




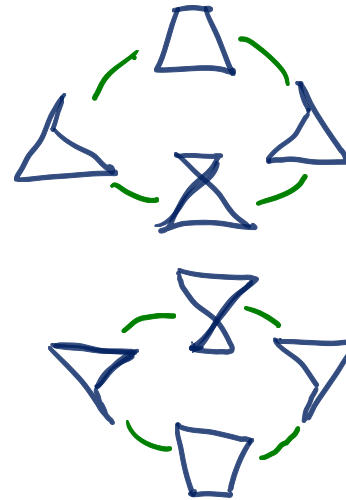
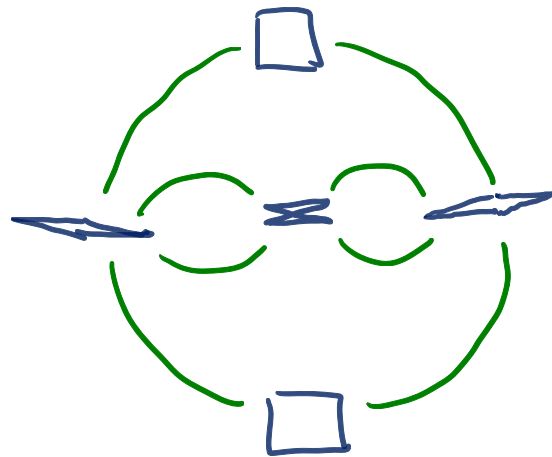
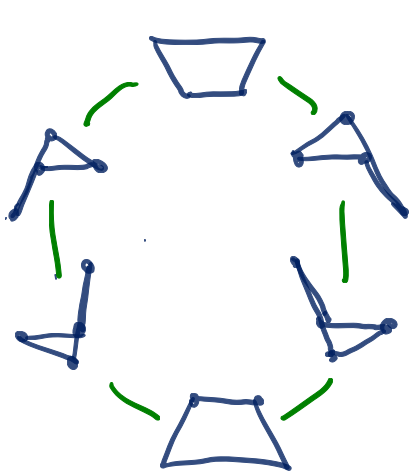
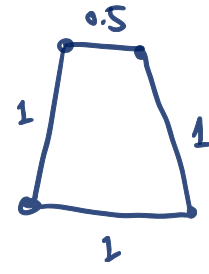
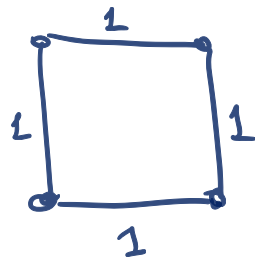
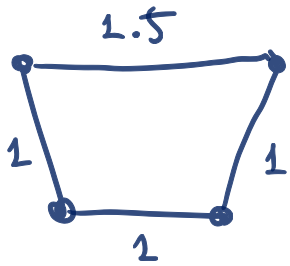


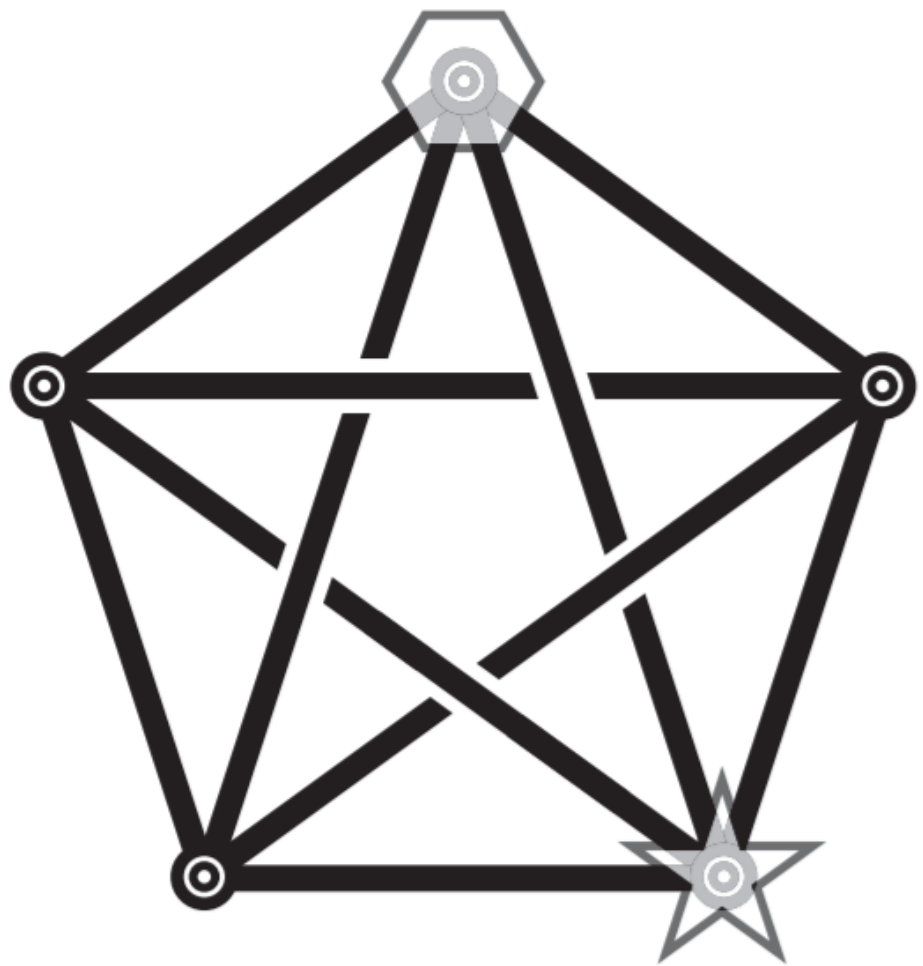
# EXAMPLE: 4-BAR LINKAGE





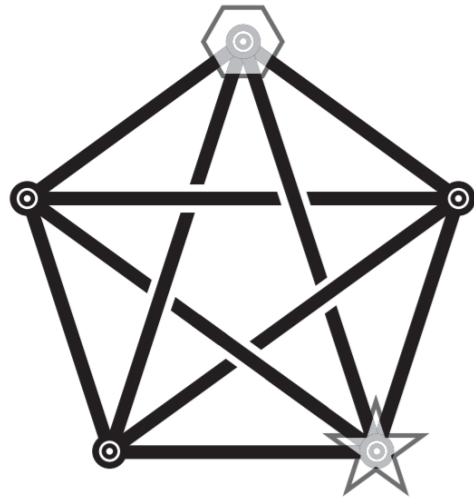
# EXAMPLE: 4-BAR LINKAGE





# EXAMPLE: SPACE OF VERTEX PAIRS



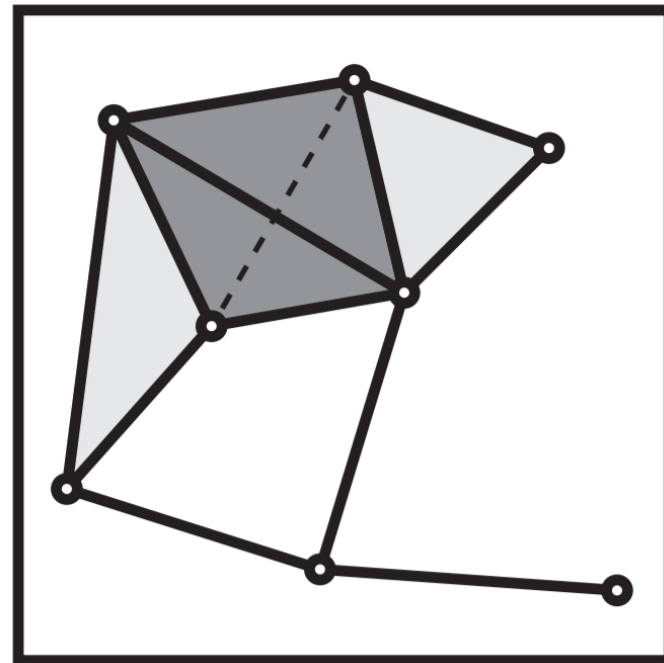
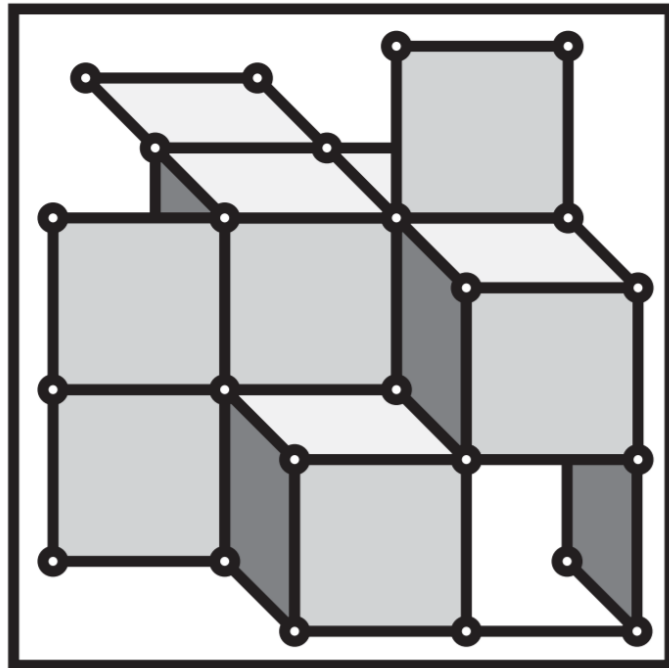


# EXAMPLE: SPACE OF VERTEX PAIRS

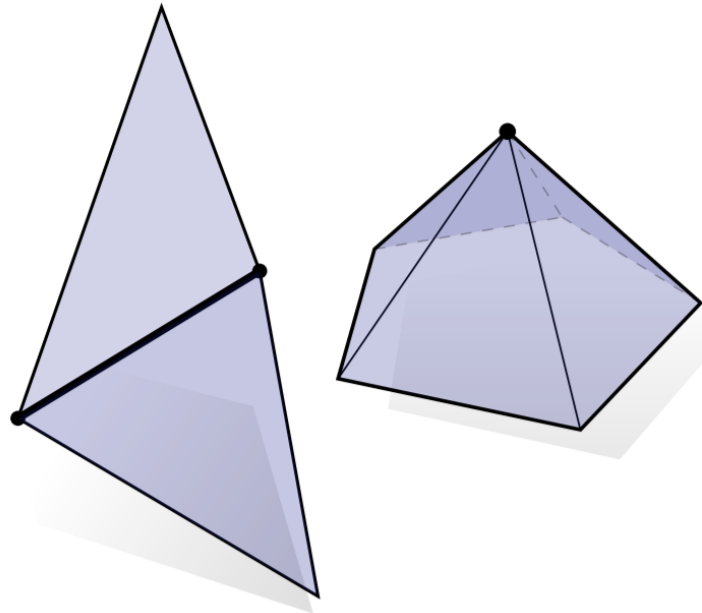


# COMPLEXES

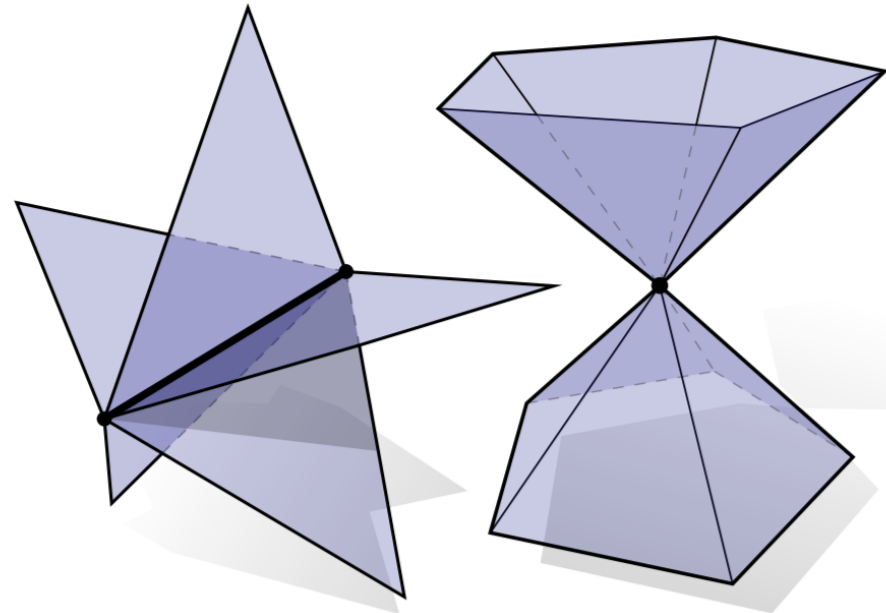
- Gluing a bunch of simplexes together



# EXTRA FEATURES

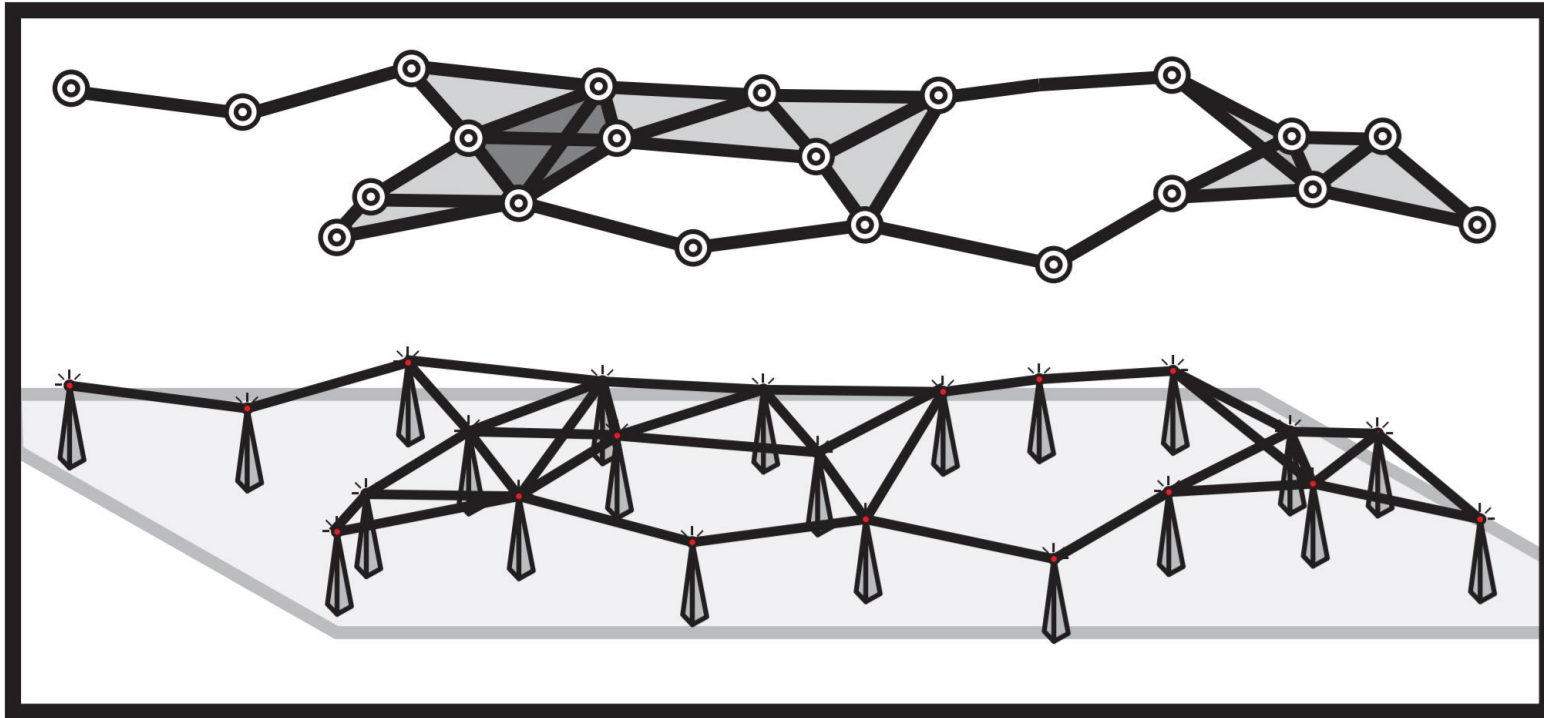


*manifold*



*nonmanifold*



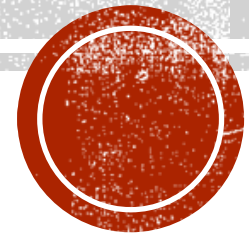


# VIETORIS-RIPS COMPLEX

- Connect any two points of distance at most  $r$
- Add all simplexes inside a clique



# **CLOSING Q. HOW DO WE TELL APART COMPLEXES?**



**CHOOSE YOUR OWN ADVENTURE:**

- (A) applications of curves and surfaces: Linkage and Folding, or**
- (B) behavior of closed curves dictates the shape of space**