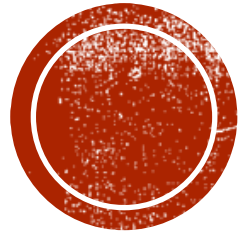


**INTRODUCTION TO  
COMPUTATIONAL  
TOPOLOGY**

**HSIEN-CHIH CHANG  
SEPTEMBER 14, 2021**



# **INTRODUCE YOURSELF!**

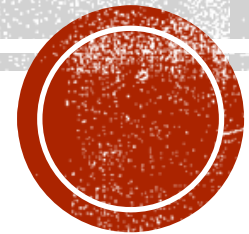
**WHO ARE YOU?**

**WHY DO YOU CARE ABOUT TOPOLOGY?**



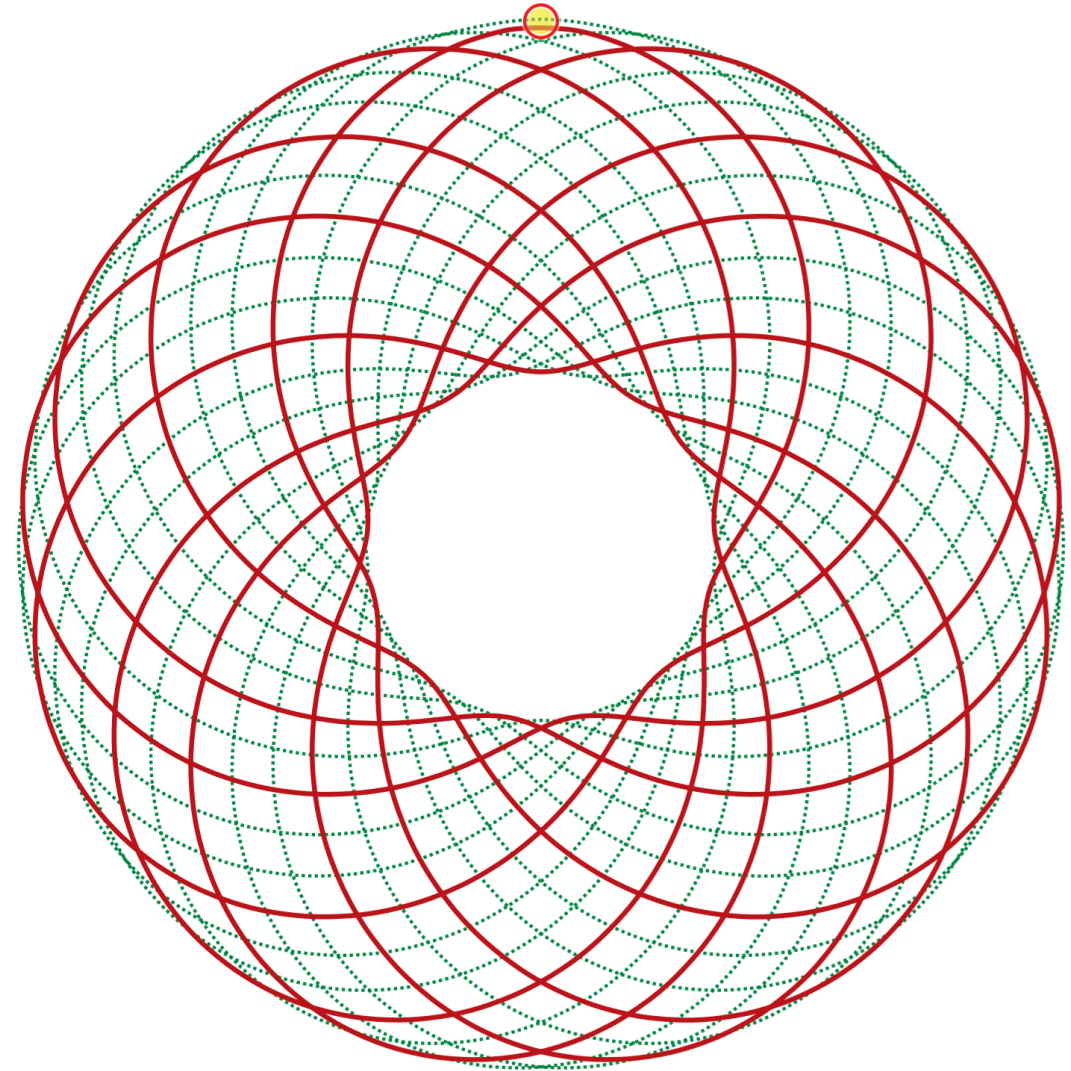
# SHAPE OF THE SPACE

## HOW DO WE COMPUTE IT?



# COURSE OUTLINE

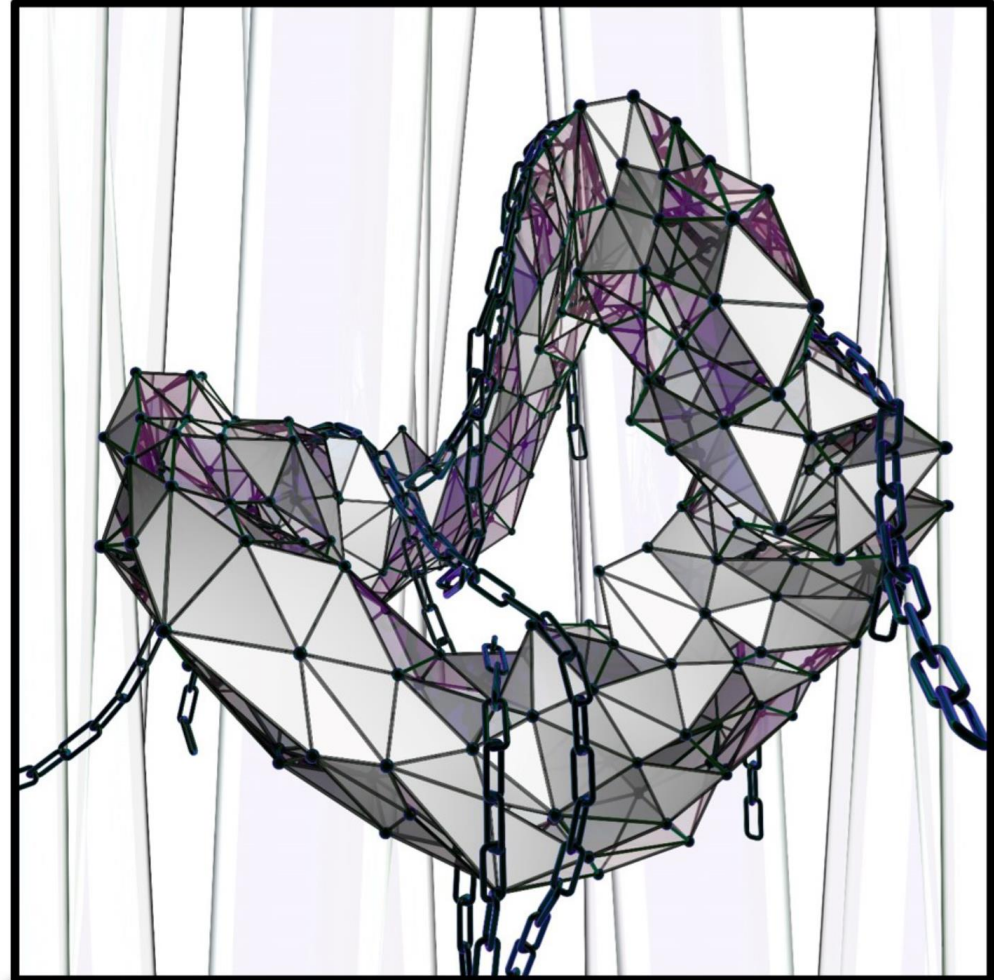
- Theory
  - Curves





# COURSE OUTLINE

- **Theory**
  - Curves
  - Surfaces/Complexes



# COURSE OUTLINE

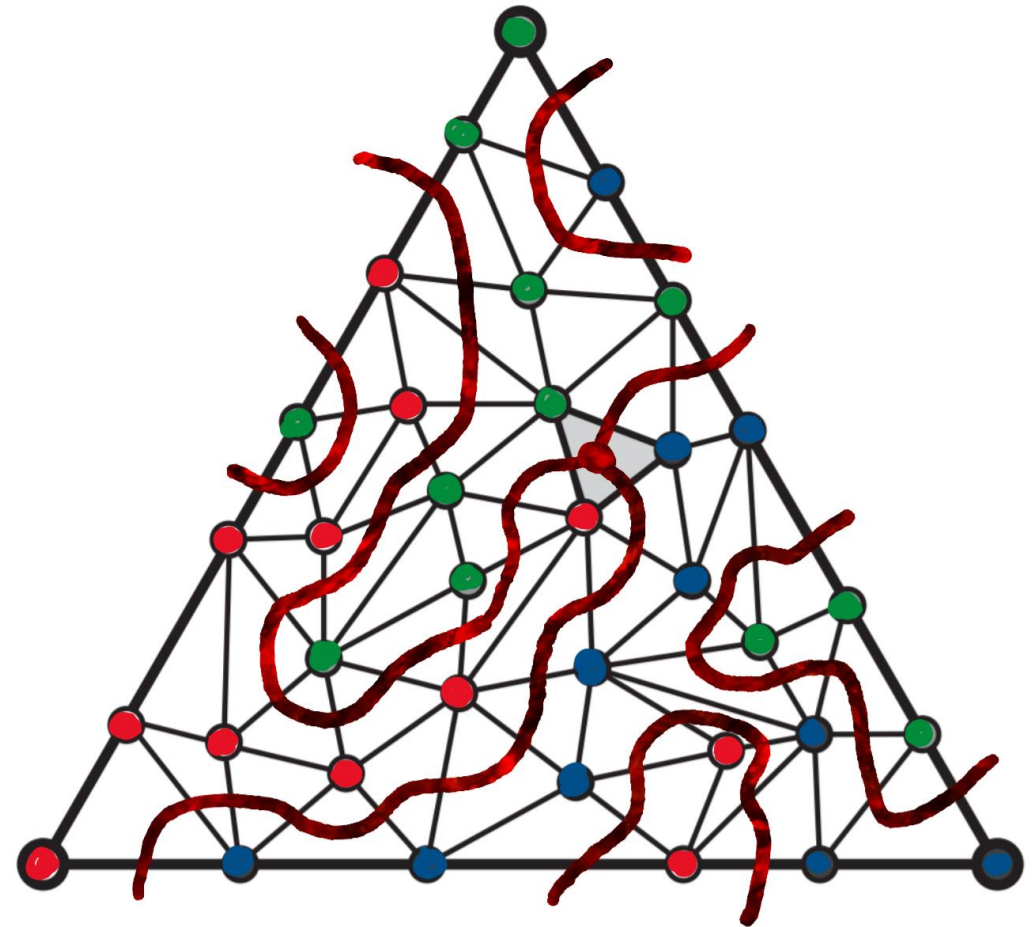
- **Theory**
  - Curves
  - Surfaces/Complexes
  - **Homotopy**



# COURSE OUTLINE

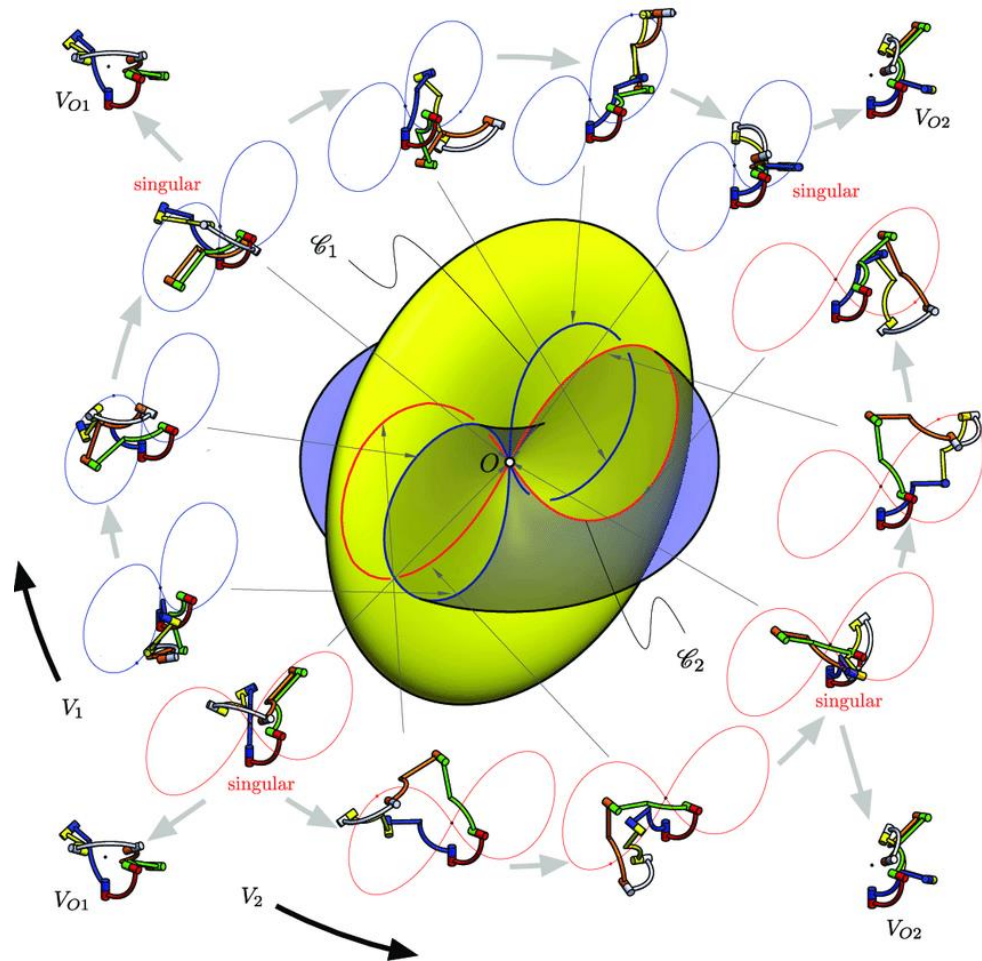
## ■ Theory

- Curves
- Surfaces/Complexes
- Homotopy
- Homology





# COURSE OUTLINE

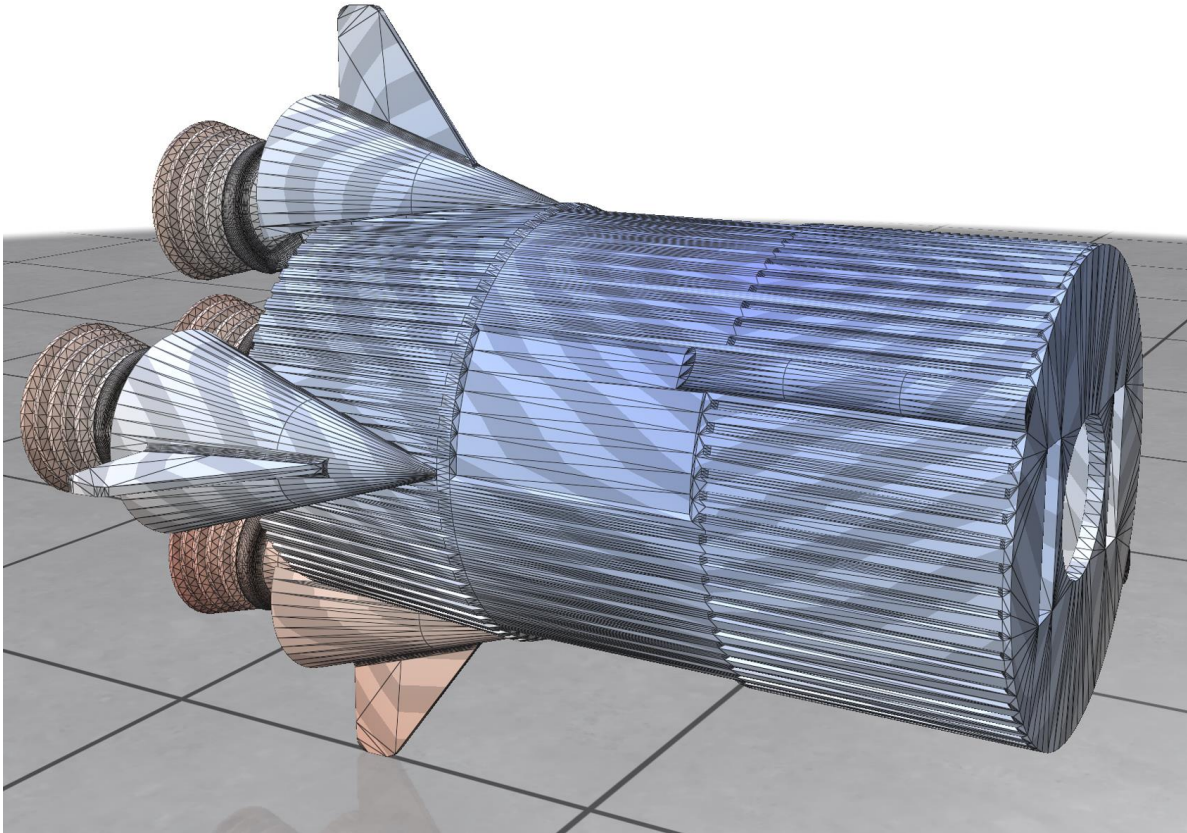


## ■ Applications

- Linkage and Folding



# COURSE OUTLINE

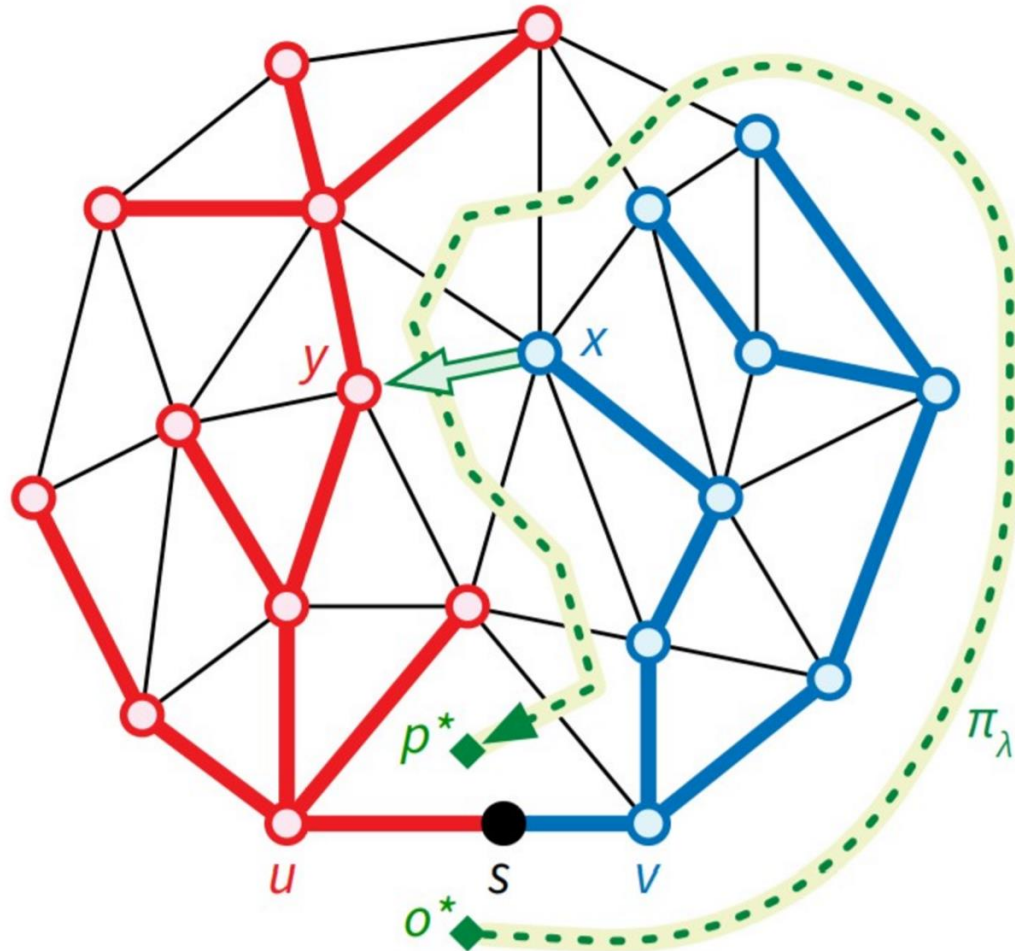


## ■ Applications

- Linkage and Folding
- Mesh Generation



# COURSE OUTLINE



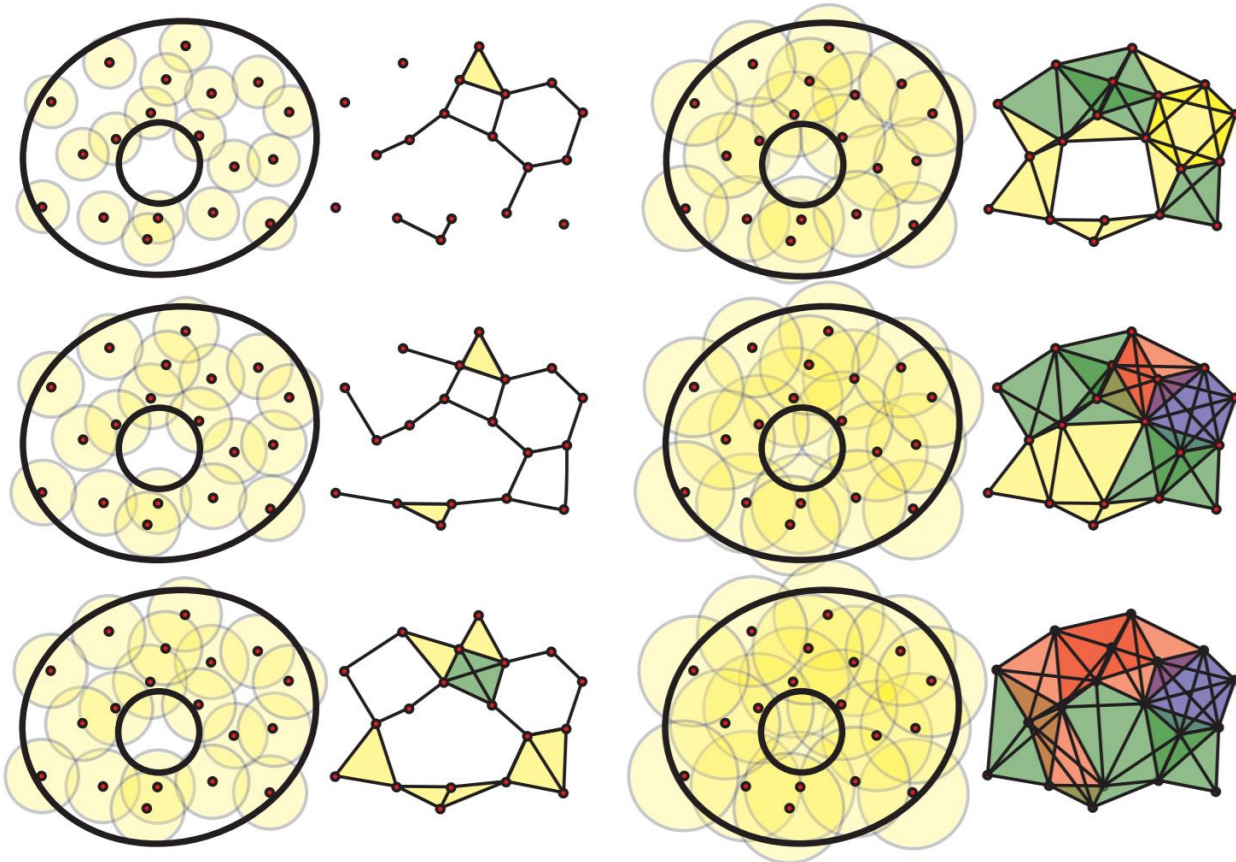
## ■ Applications

- Linkage and Folding
- Mesh Generation
- Optimization





# COURSE OUTLINE

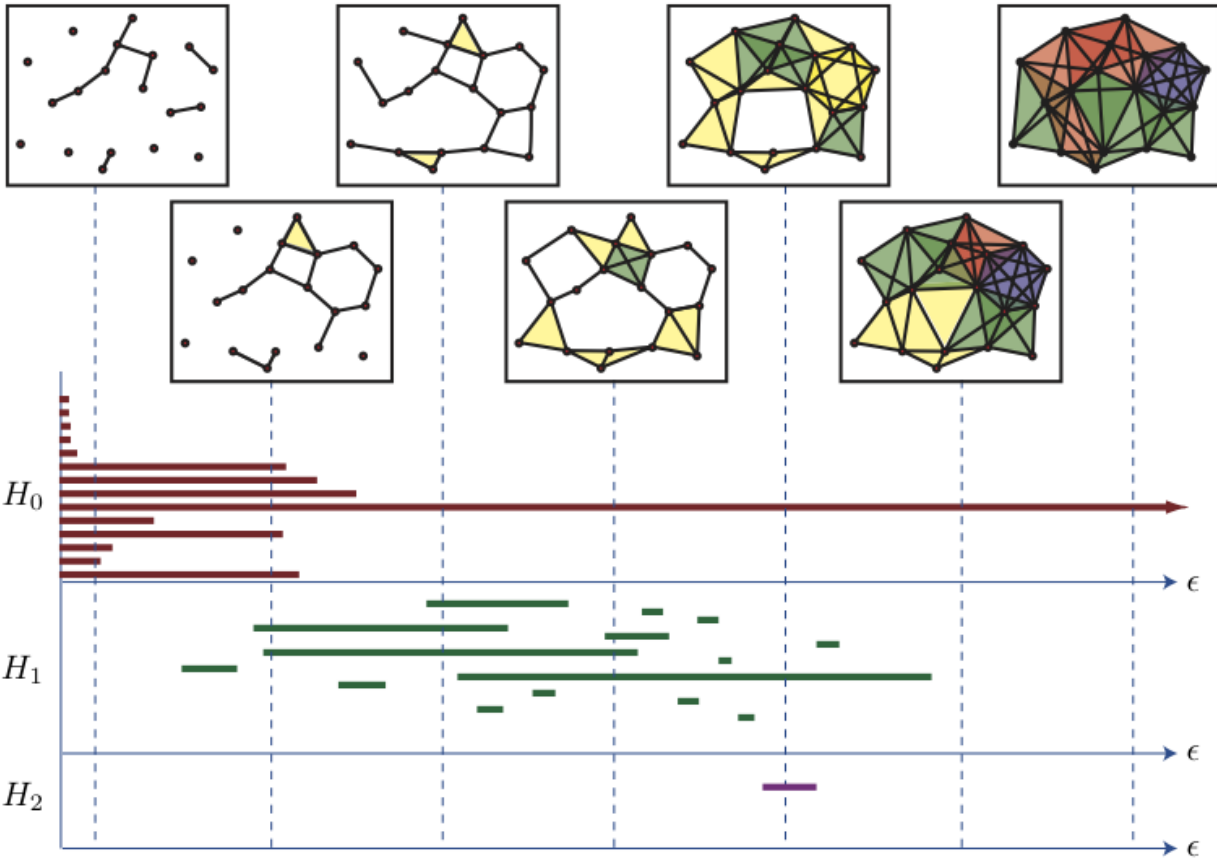


## ■ Applications

- Linkage and Folding
- Mesh Generation
- Optimization
- Topological Data Analysis



# COURSE OUTLINE



## ■ Applications

- Linkage and Folding
- Mesh Generation
- Optimization
- **Topological Data Analysis**



# COURSE OUTLINE

## ■ Theory

- Curves
- Surfaces/Complexes
- Homotopy
- Homology

## ■ Applications

- Linkage and Folding
- Mesh Generation
- Optimization
- Topological Data Analysis



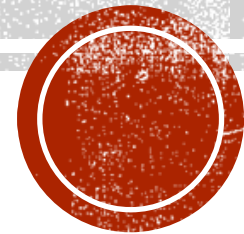


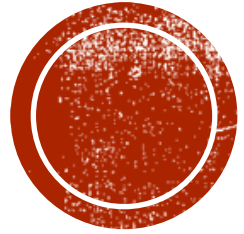
# LOGISTICS

- **Course webpage**
- **Lectures**
- **Office Hours**
- **Discussion forum**
- **Homework**
- **Project**



**STOP ME IF YOU ARE LOST**





# **SIMPLE PLANAR CURVES**

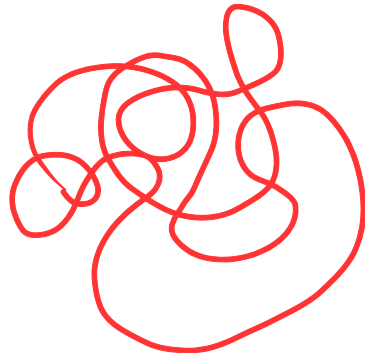
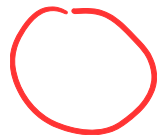




# FORMAL DEFINITIONS

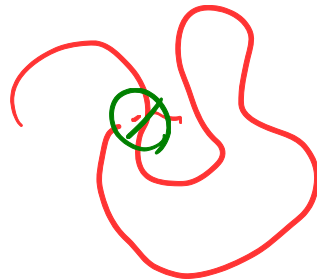
- Planar curve

$$\gamma: S^1 \rightarrow \mathbb{R}^2$$



- Simple

$\gamma$  embedding  
injective



no self-crossings



# SIMPLE PLANAR CURVES

Jordan curve

[Jordan 1887]

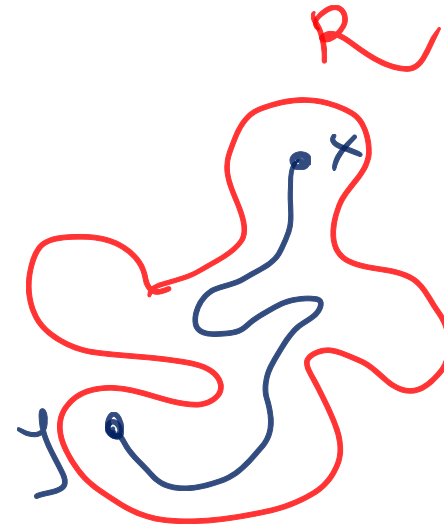


# FORMAL DEFINITIONS

- **Connected**

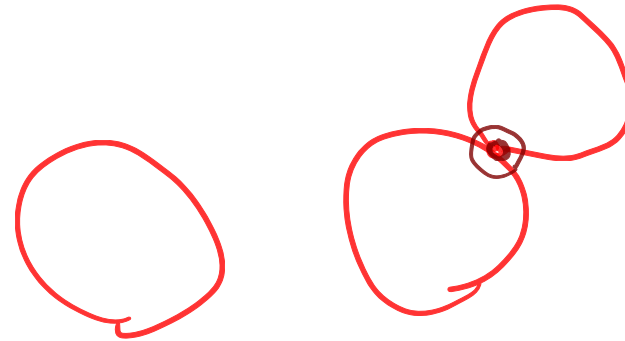
$x \sim y$  if  $\exists \pi: [0,1] \rightarrow \mathcal{R}$ .  
 $\pi(0) = x, \pi(1) = y$

$\mathcal{R} \subseteq \mathbb{R}^2$   
 $\mathcal{R}$  connected if  $\forall x, y \in \mathcal{R}, x \sim y$

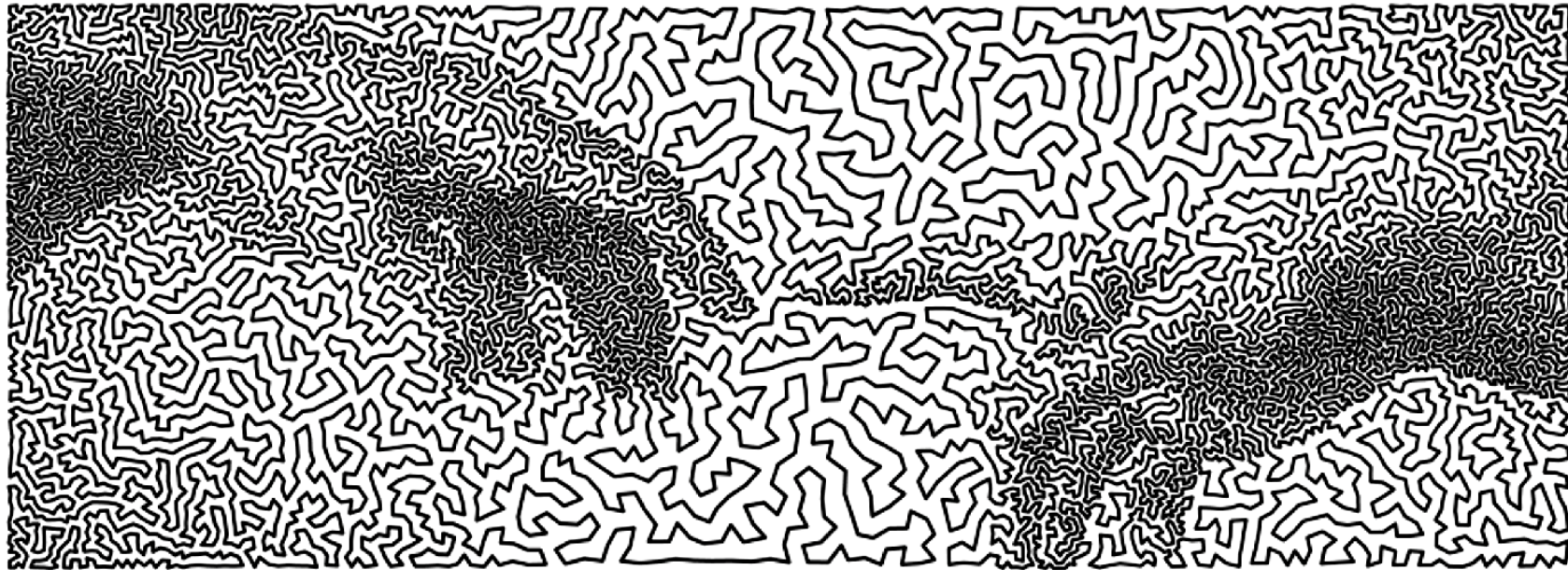


- **Connected component**

maximal subregions that are connected.



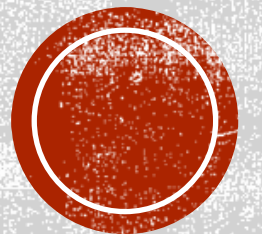




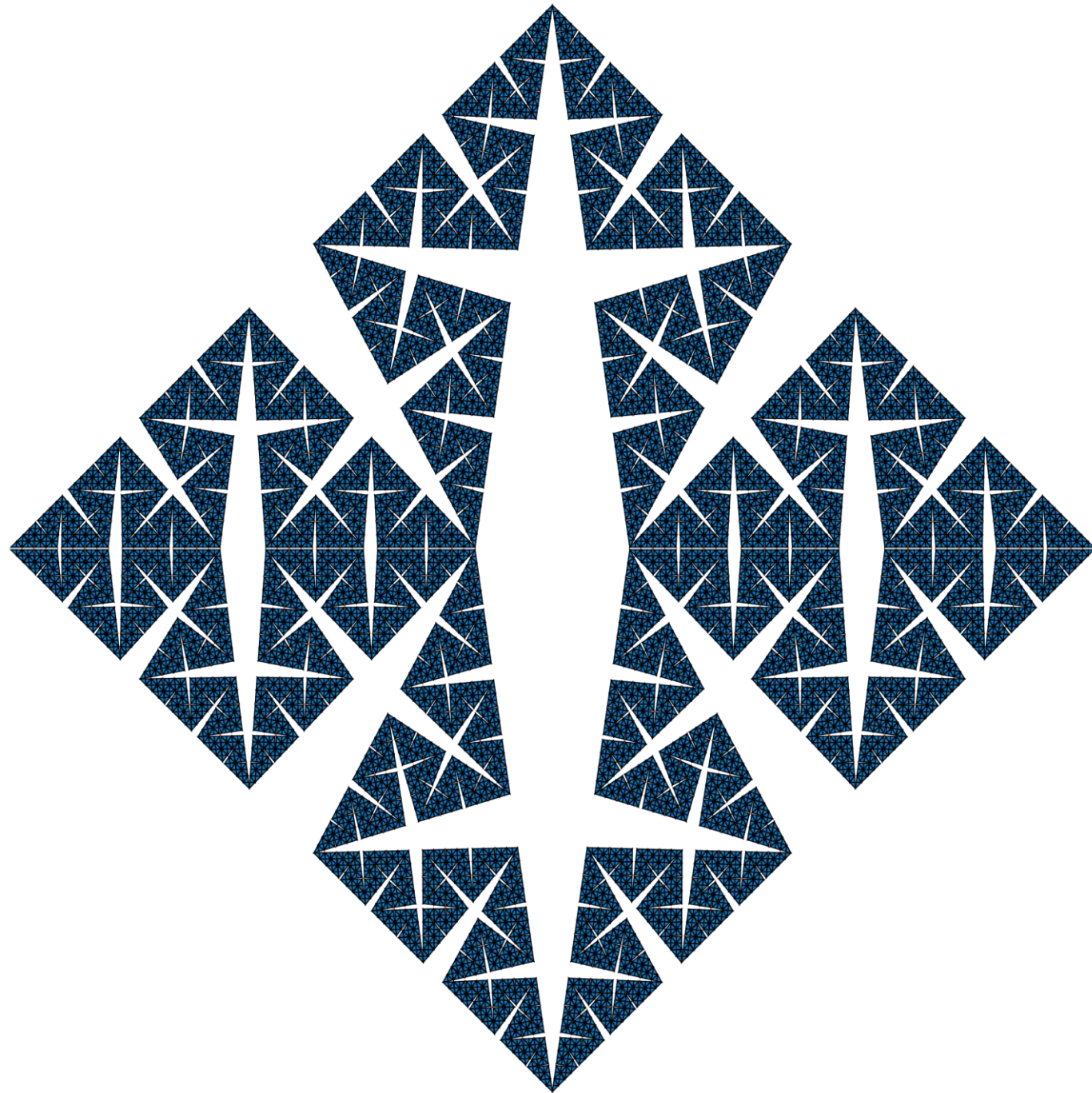
# JORDAN CURVE THEOREM

[Bolzano ~1800s] [Jordan 1887]

Any simple closed curve  $C$  separates  $\mathbb{R}^2 \setminus C$  into exactly two connected components







# SIMPLE PLANAR CURVES

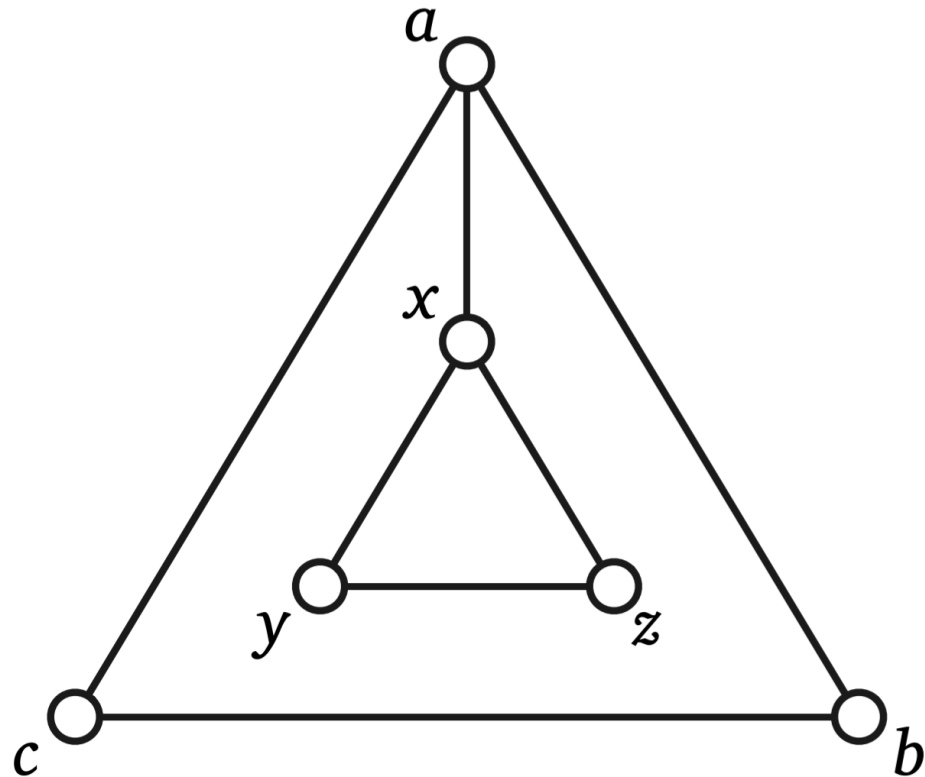
Osgood curve

[Osgood 1903]

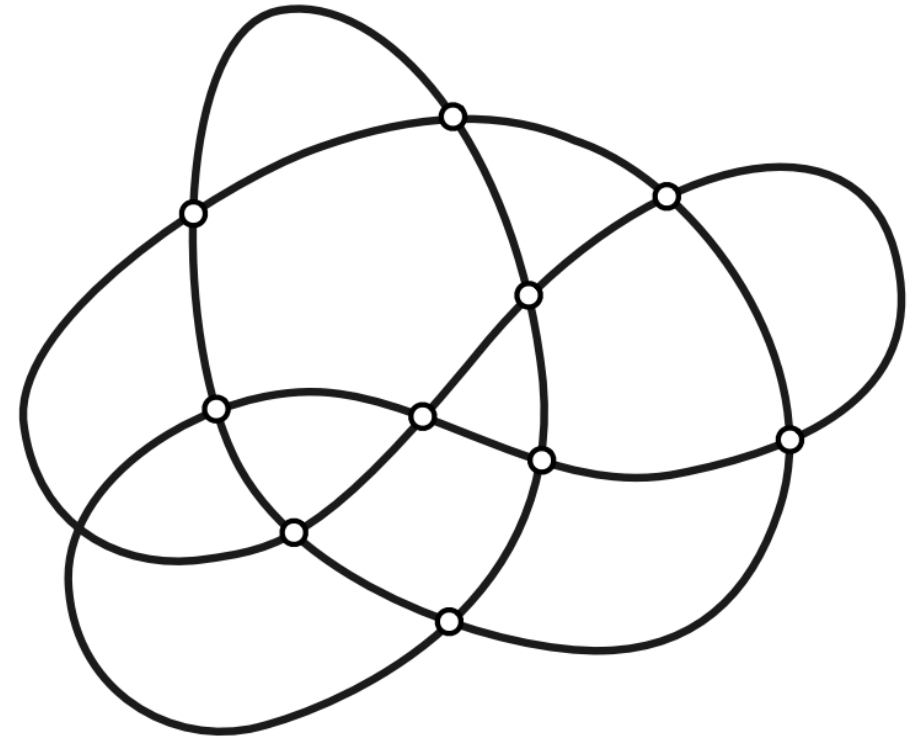


# REPRESENTATION OF CURVES

## ■ Polygonal

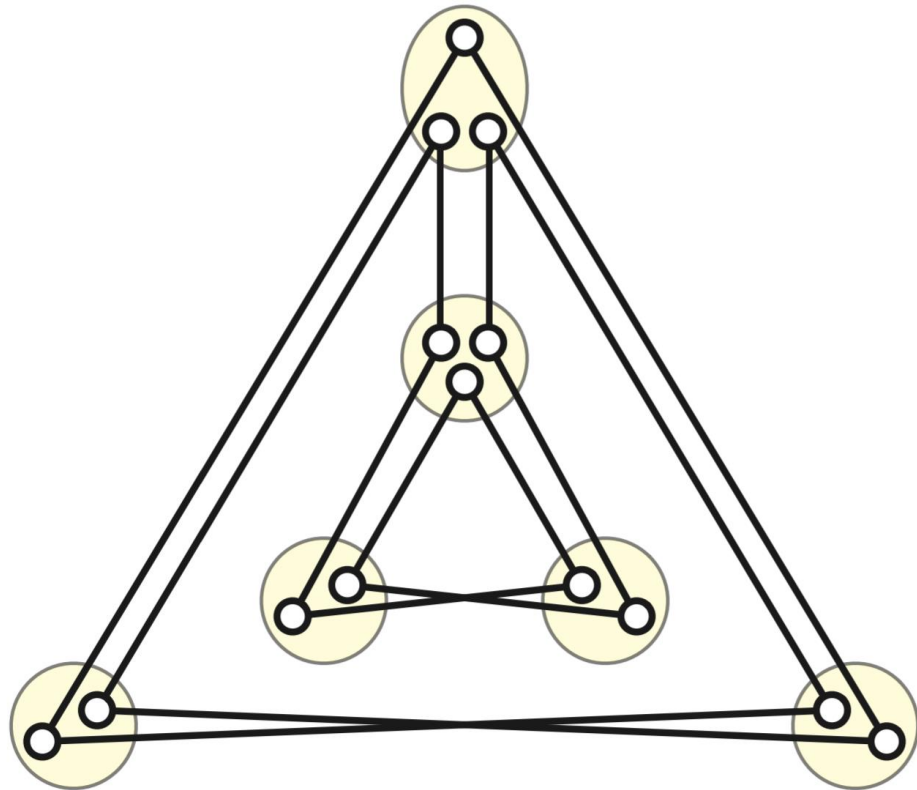


## ■ Generic

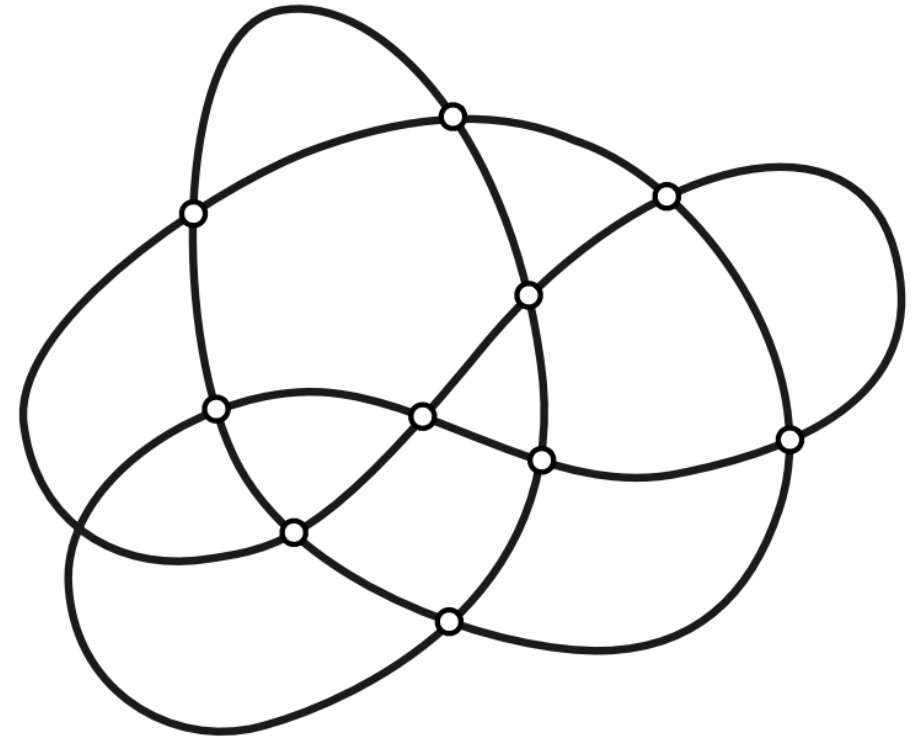


# REPRESENTATION OF CURVES

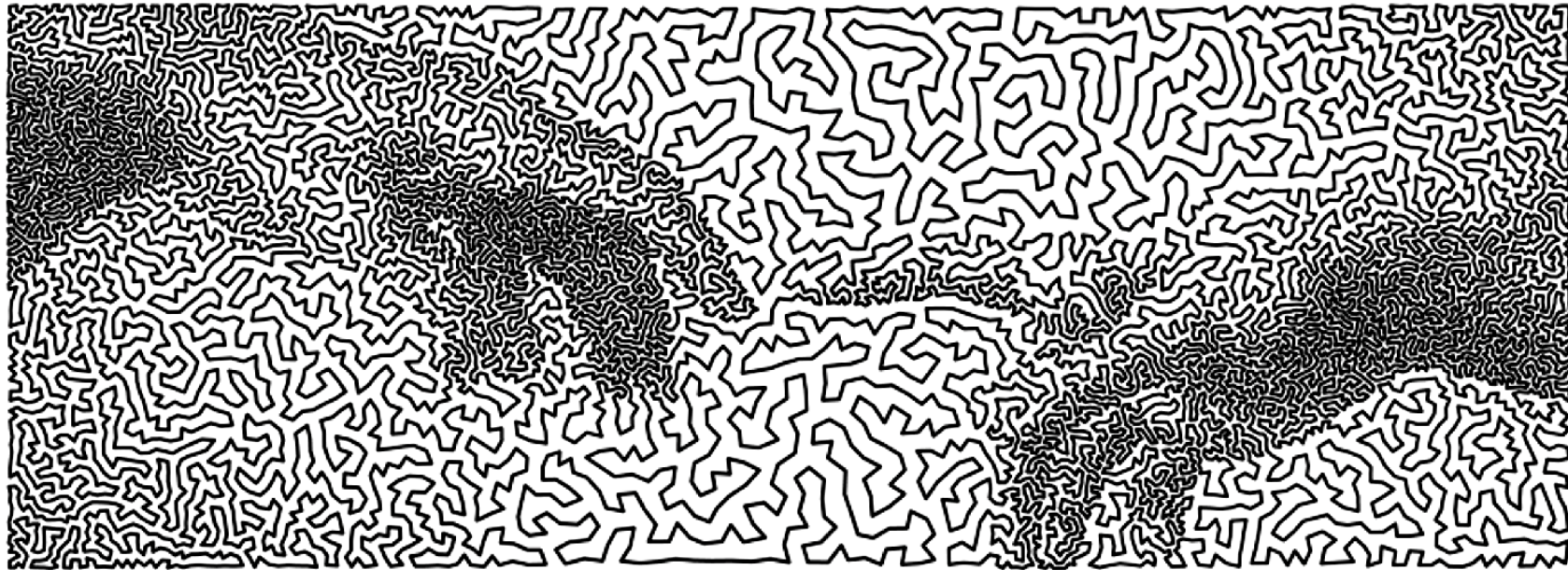
■ Polygonal



■ Generic

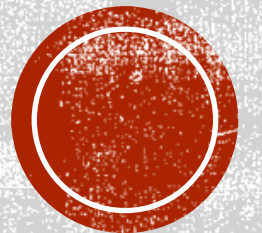






# JORDAN POLYGON THEOREM [Dehn 1899]

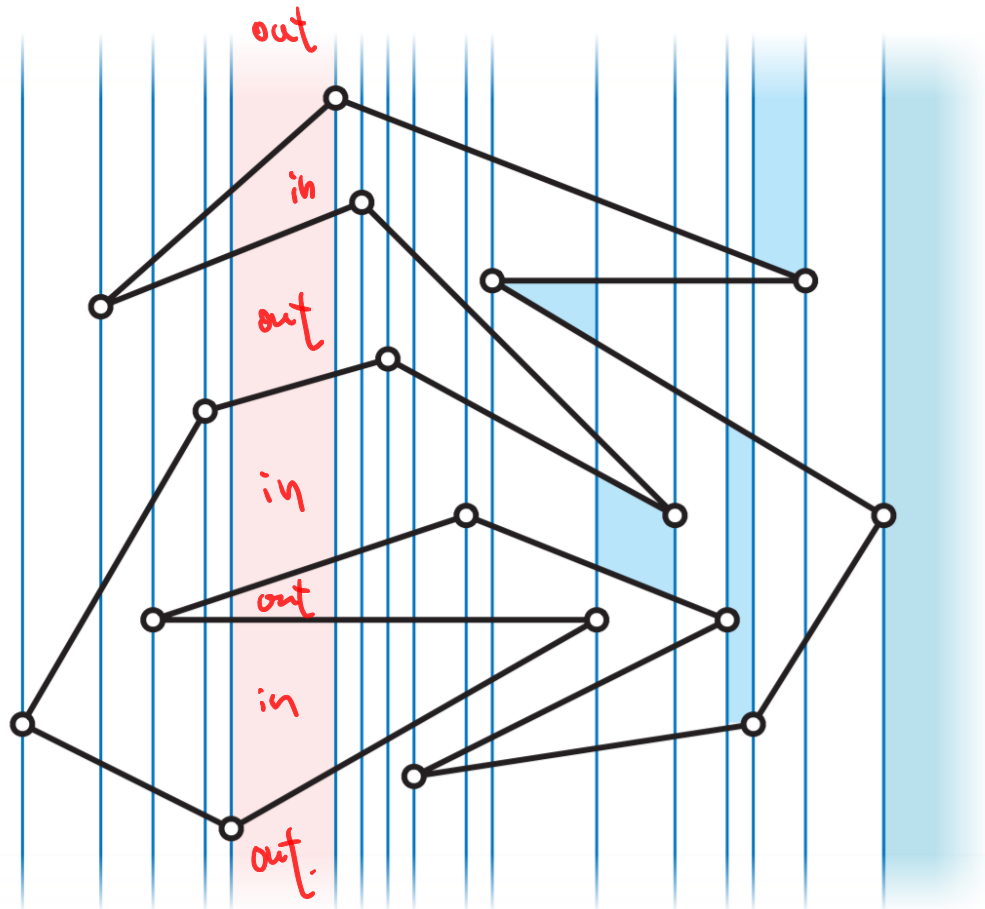
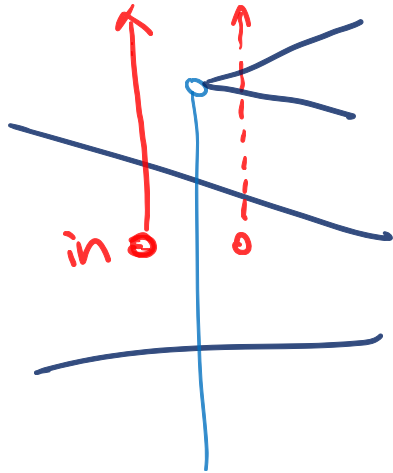
Any simple **polygon**  $P$  separates  $\mathbb{R}^2 \setminus P$  into exactly two connected components





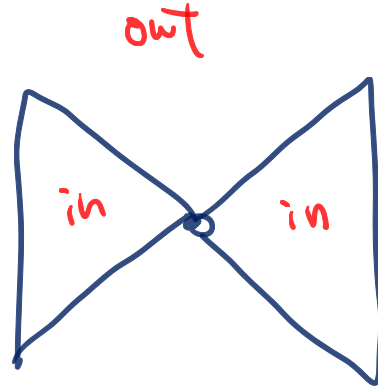
# PROOF OF JORDAN POLYGON THEOREM

- Intuition: Parity argument

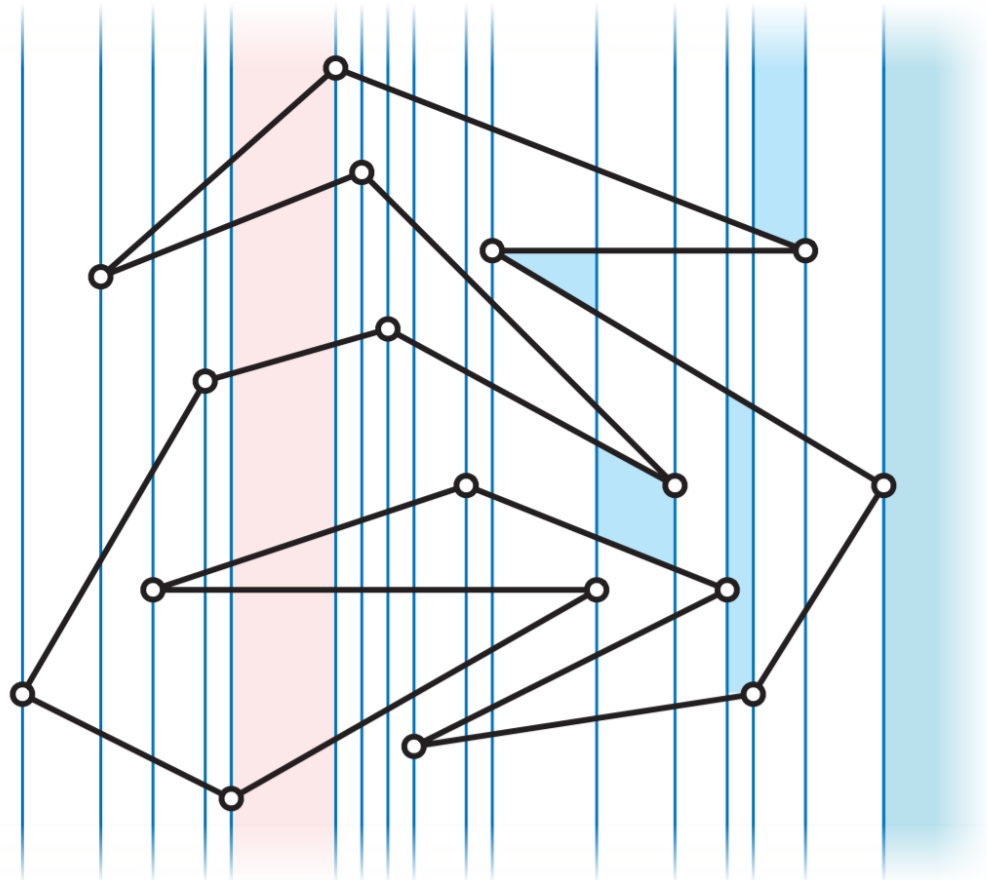


# PROOF OF JORDAN POLYGON THEOREM

- Lemma  $\geq 2$ 
  - Parity argument

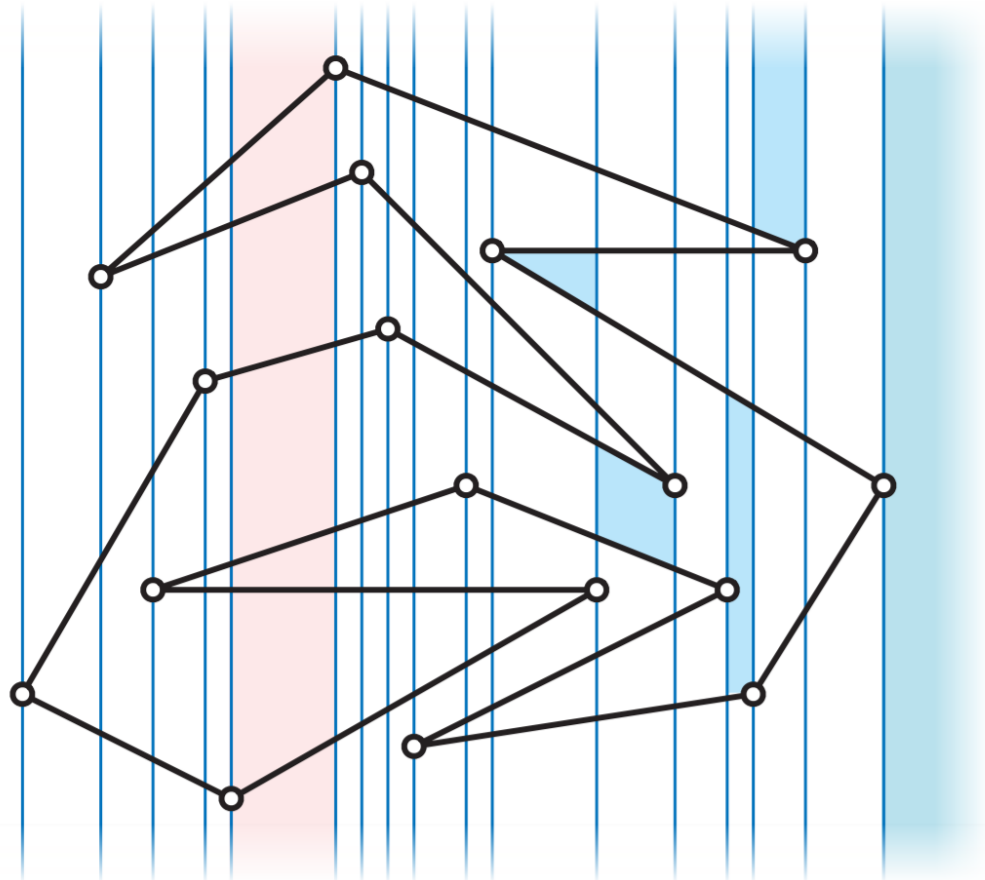
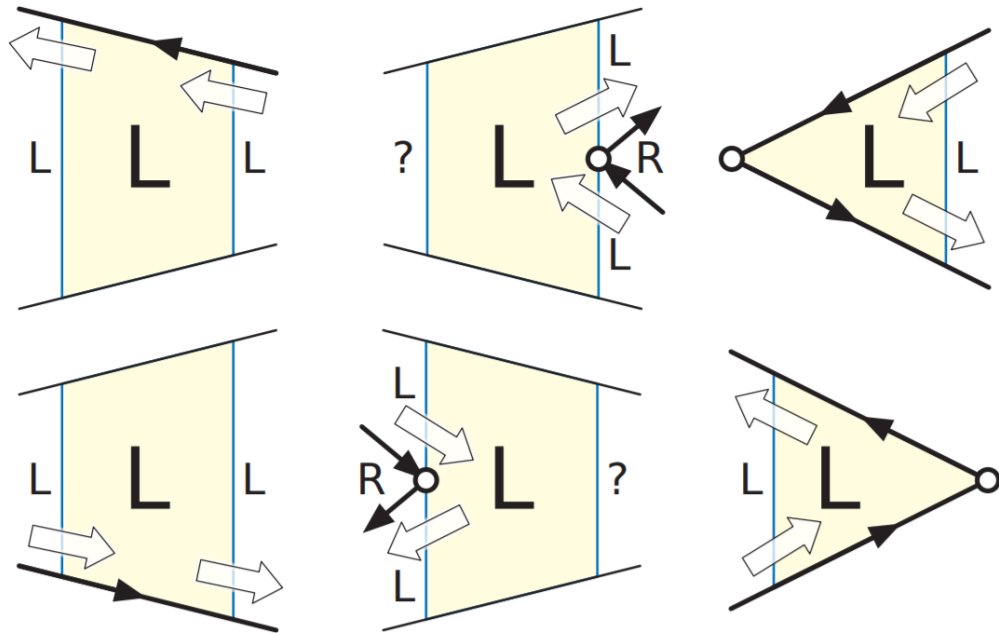


- Lemma  $\leq 2$ ?



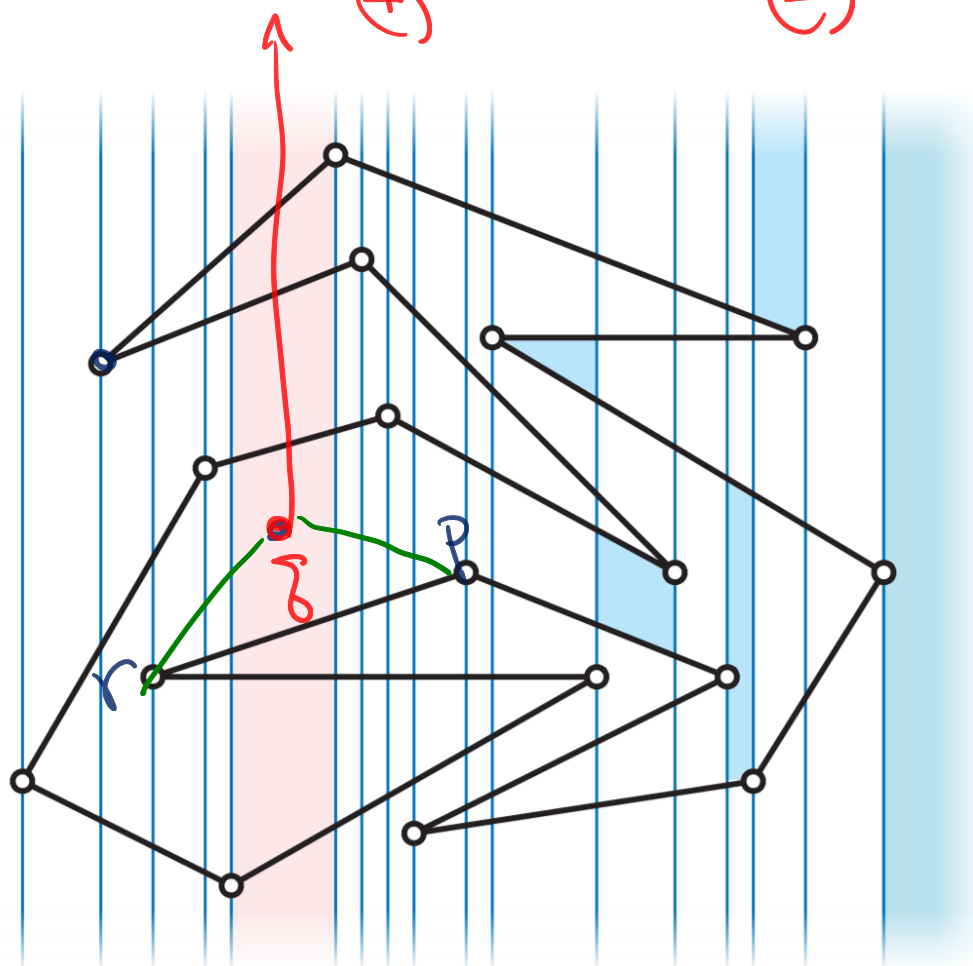
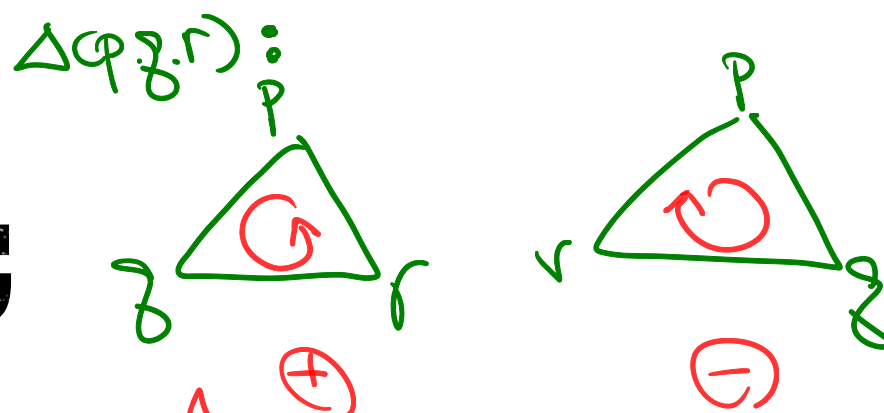
# PROOF OF JORDAN POLYGON THEOREM

## ■ Lemma $\leq 2$



$$\Delta(p.g.r) = \begin{vmatrix} 1 & p.x & p.y \\ 1 & g.x & g.y \\ 1 & r.x & r.y \end{vmatrix}$$

# INSIDE-POLYGON TESTING



Inside Polygon? (P-g) :

sign  $\leftarrow 0$

for each segment  $\vec{pr}$  :

$\Delta \leftarrow \Delta(p.g.r)$

sign of triangle  
(p.g.r)

if  $p.x \leq g.x < r.x$  :

sign  $\leftarrow -\Delta \cdot \text{sign}$

crossing  
above  
red

if  $r.x \leq g.x < p.x$  :

sign  $\leftarrow \Delta \cdot \text{sign}$

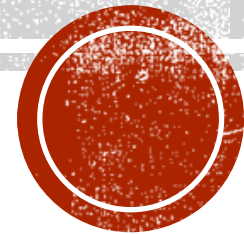
crossing  
below  
red.

return sign



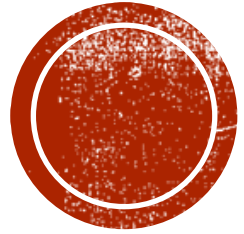


# INTERMISSION



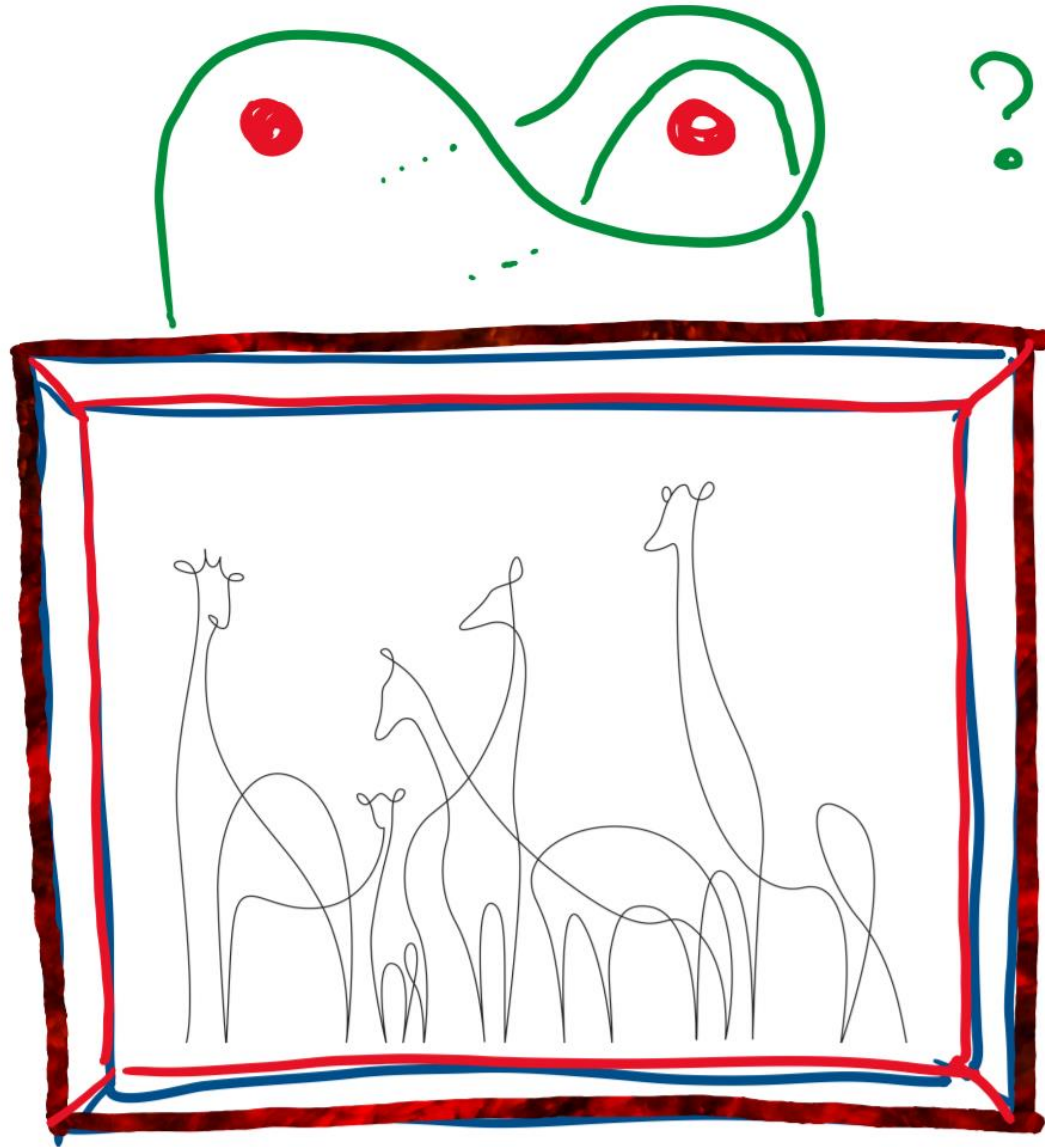
**FOOD FOR THOUGHT.**

**How to compute the area of a simple polygon?**



# WINDING NUMBER



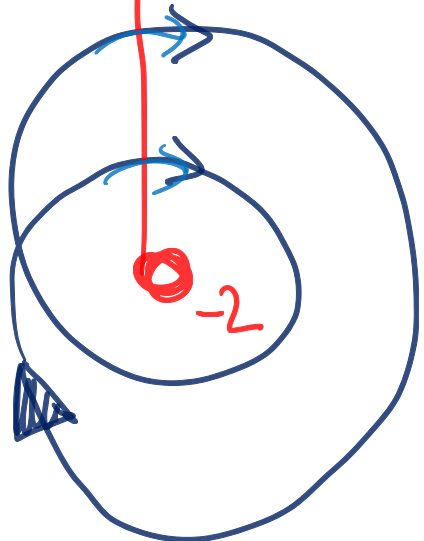
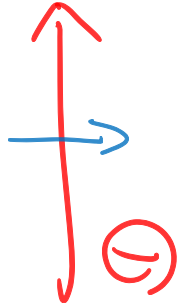
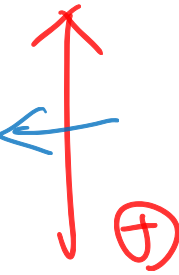
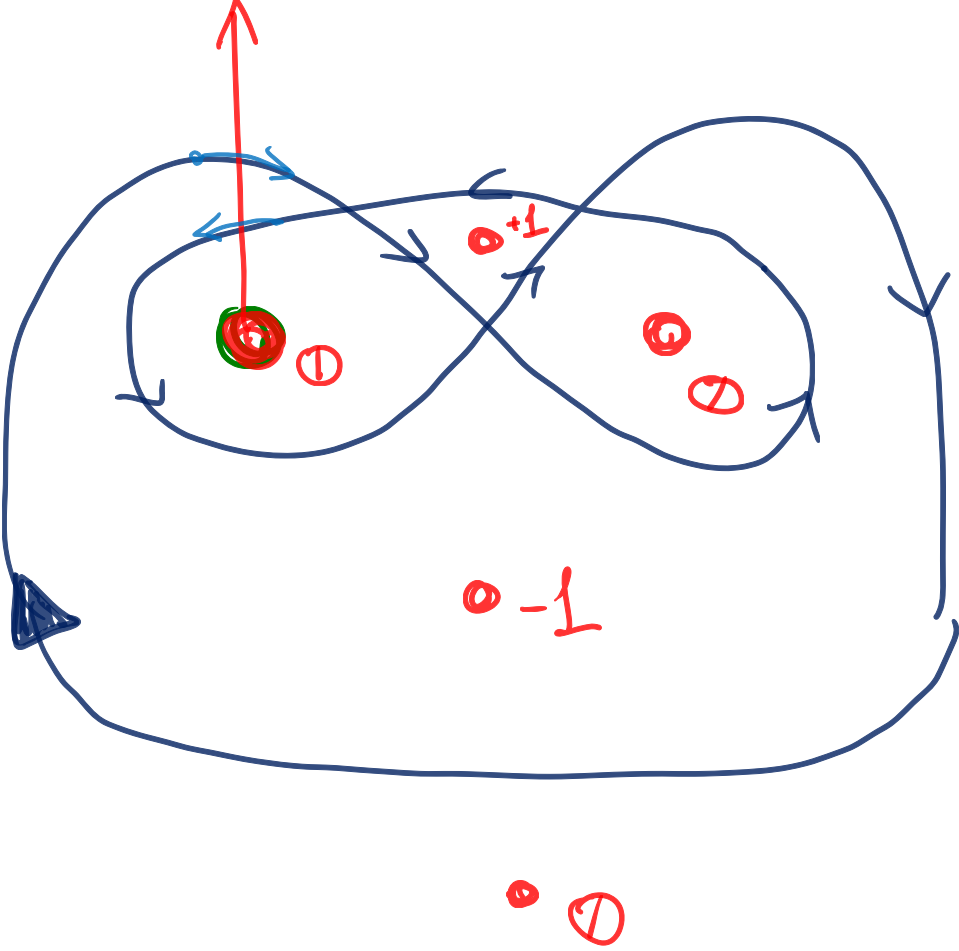


# PICTURE-HANGING PUZZLE

Can you hang a picture  
with two nails, such that  
the picture falls if  
**either** nail is pulled?



# LOOK AT THE SOLUTION





# COMPUTING WINDING NUMBER

Winding Number (P,  $\gamma$ ) :

wind  $\leftarrow 0$

for each segment  $\overrightarrow{pr}$  :

$\Delta \leftarrow \Delta(p, \gamma, r)$

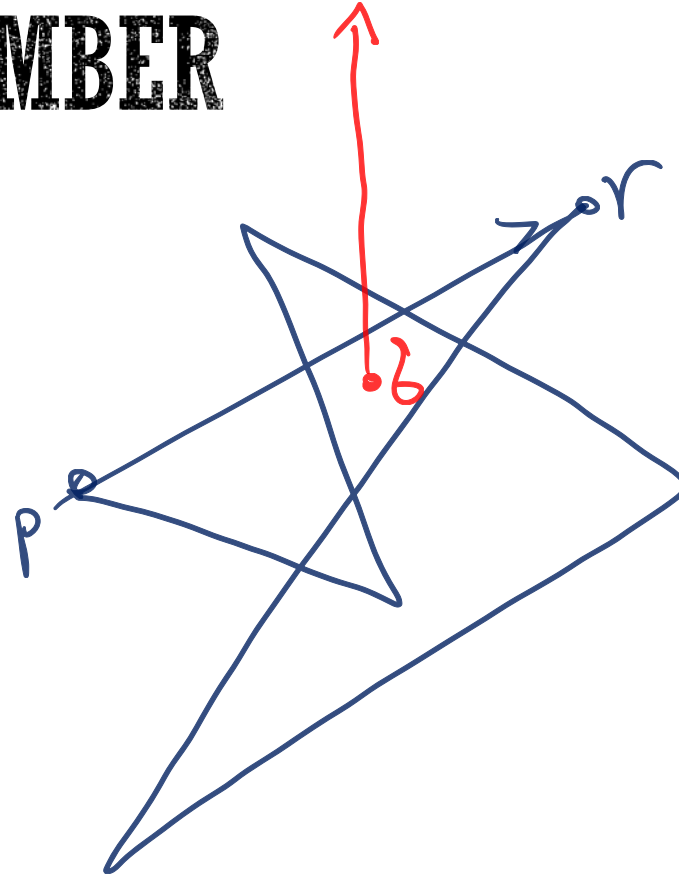
if  $p.x \leq \gamma.x < r.x$   
and  $\Delta = +1$  :

wind --

if  $r.x \leq \gamma.x < p.x$   
and  $\Delta = -1$  :

wind ++

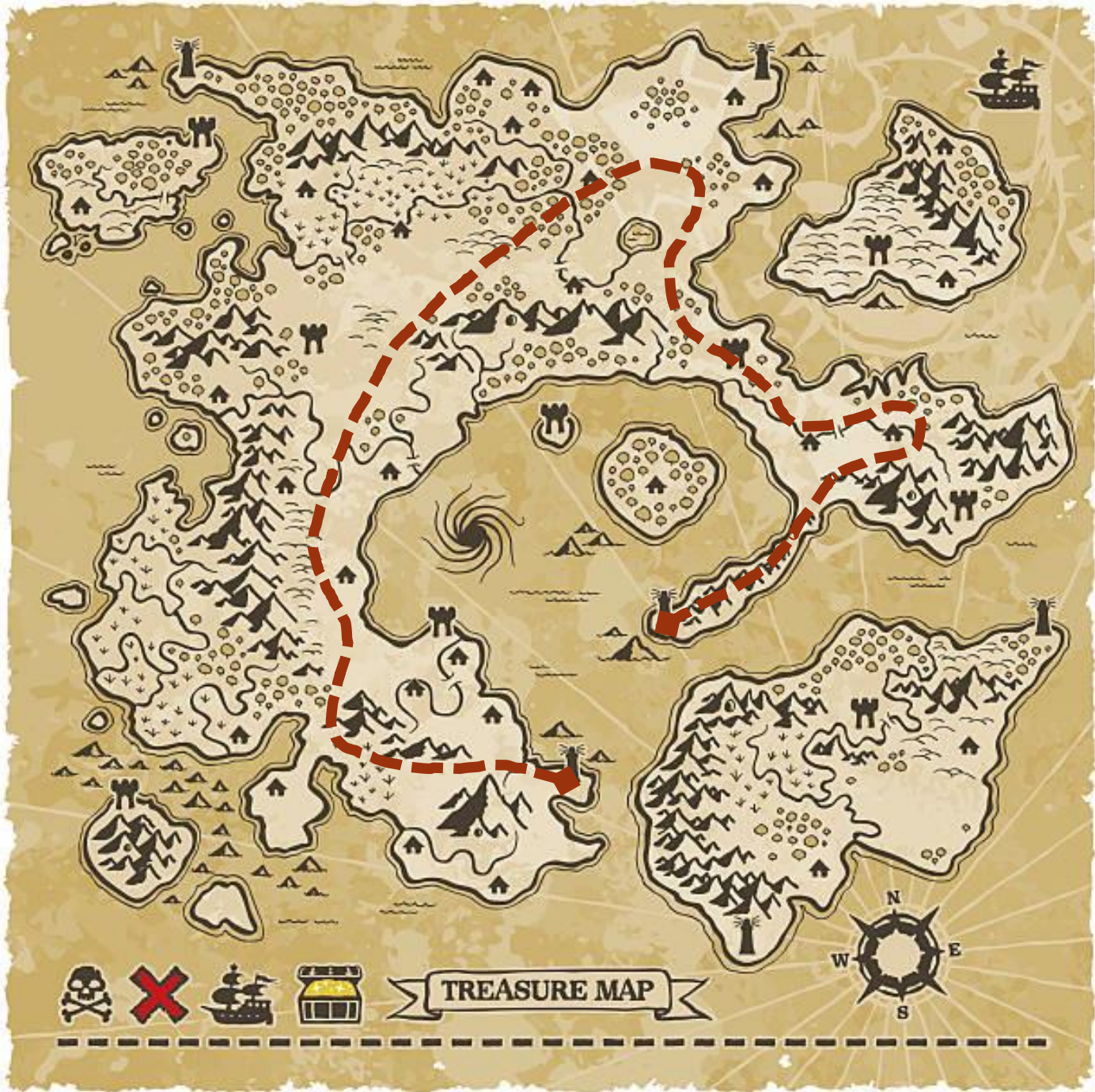
return wind



# **WIND<sub>q</sub>(P) INVARIANT UNDER MORPHING**

- **Homotopy**

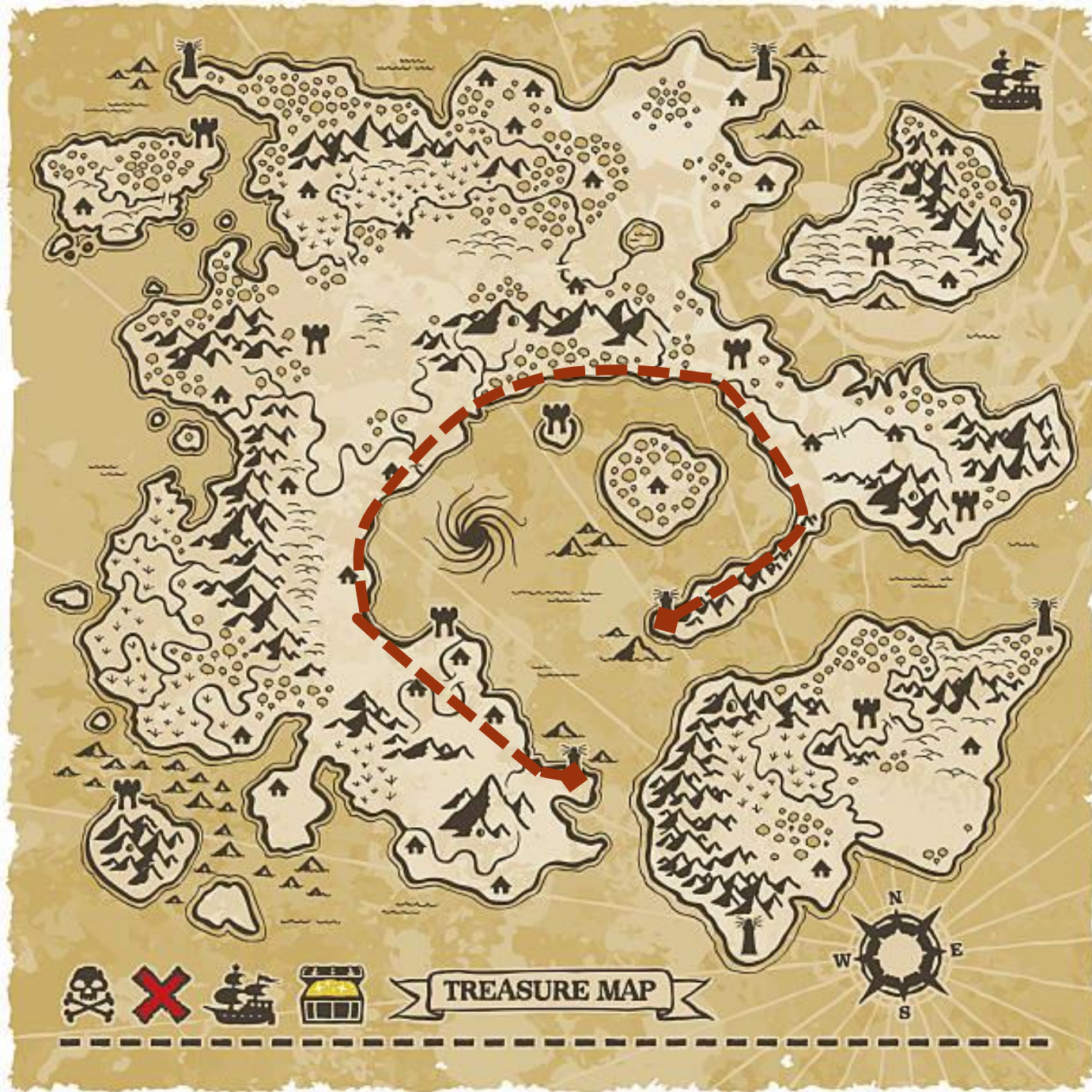




# CURVE MORPHING







# CURVE MORPHING





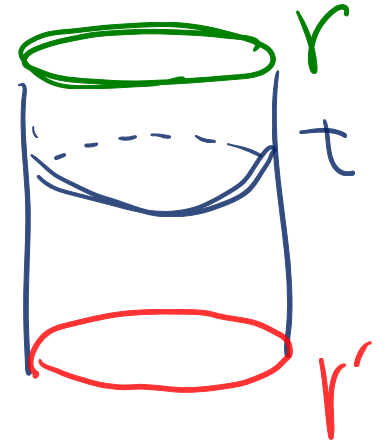
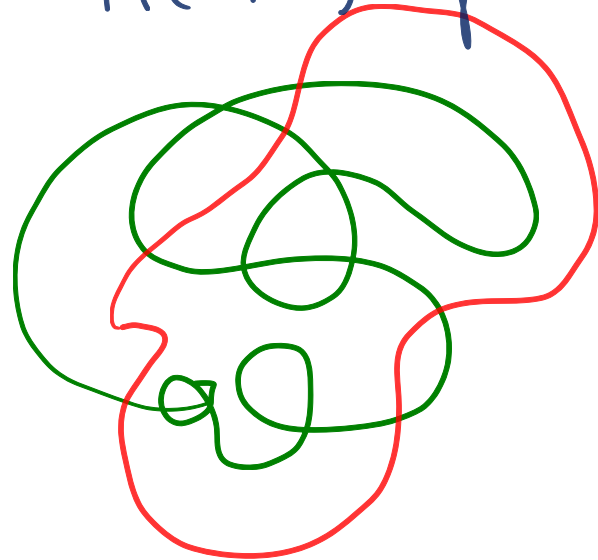
# WIND<sub>q</sub>(P) INVARIANT UNDER MORPHING

## ■ Homotopy

cont. change from  $\gamma$  to  $\gamma'$

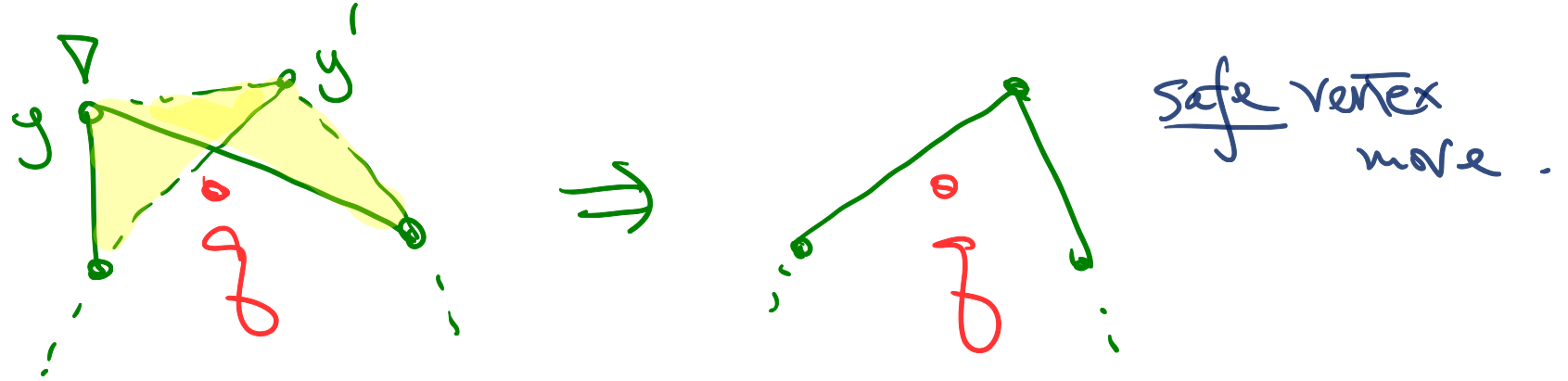
$$H: S^1 \times [0, 1] \rightarrow \mathbb{R}^2$$

$$H(\cdot, 0) = \gamma \quad H(\cdot, 1) = \gamma'$$

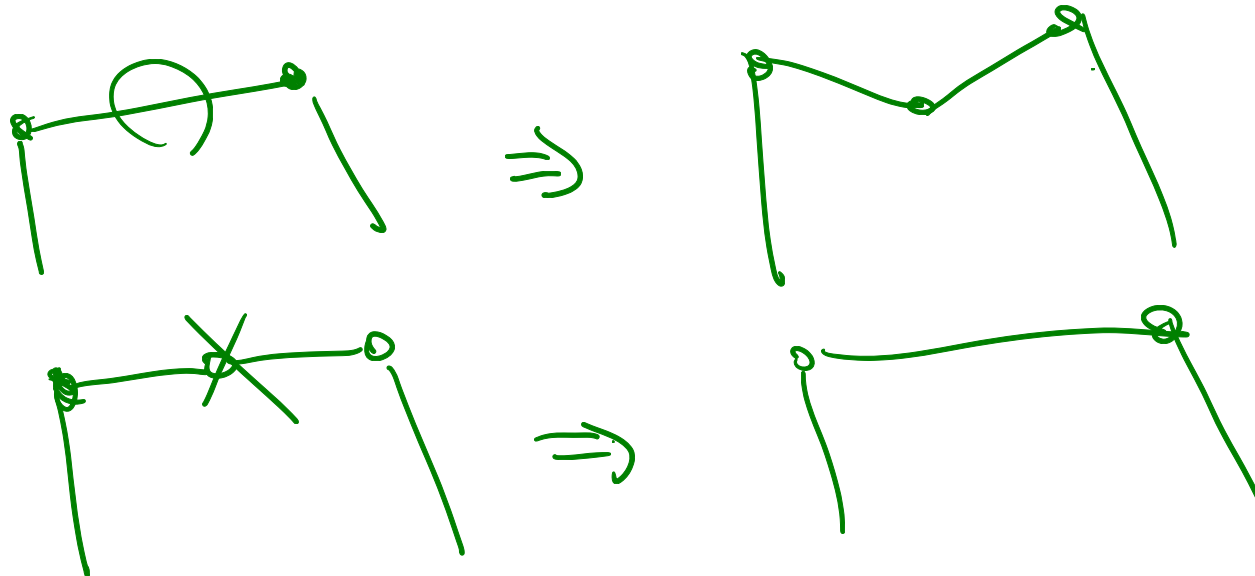


# $\text{WIND}_q(P)$ INVARIANT UNDER MORPHING

■ Homotopy

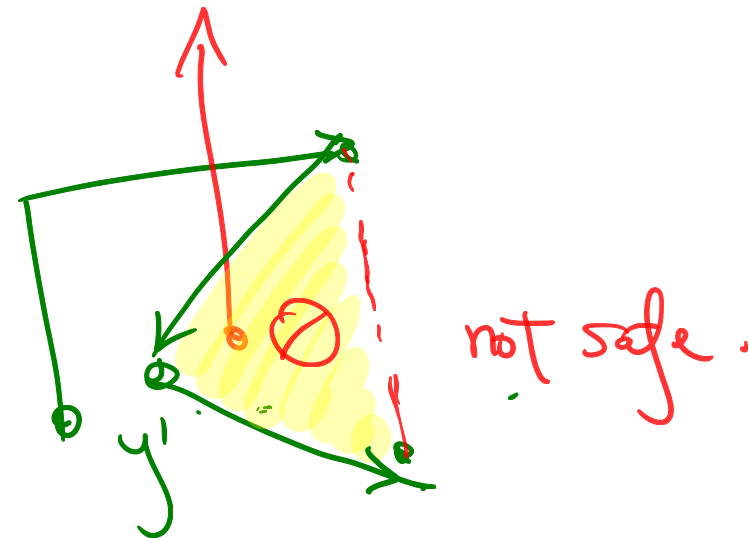
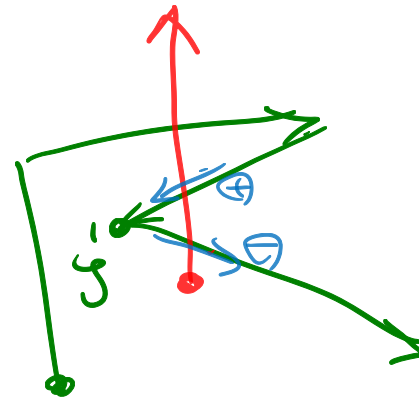
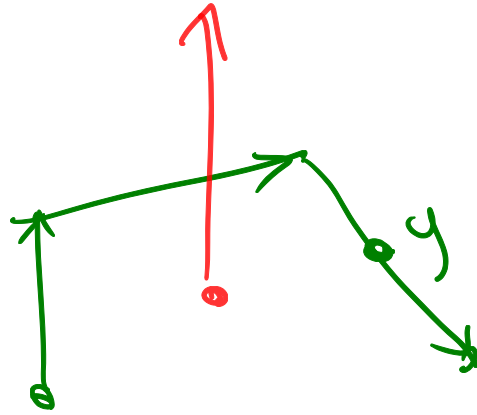


■ Vertex move



**THEOREM.**  $\text{Wind}_q(P)$  is invariant under safe vertex moves.

pf sketch.

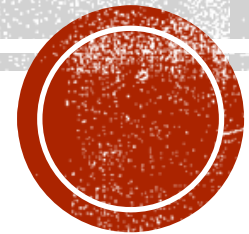


**THEOREM.** Two polygons  $P$  and  $Q$  are homotopic in  $\mathbb{R}^2 \setminus q$  if and only if they have the same  $\text{Wind}_q$ . [Hopf 1935]





# **WIND<sub>q</sub> IS A COMPLETE HOMOTOPIC INVARIANT!**



## **TAKEAWAY.**

**Planar curve can be described by how many times it goes around reference points.**