

1. **Busy chef.** Construct NFAs that recognize the following languages.

(a) Let *Sandwich* be an automatic language.

$$\text{Cut}(\text{Sandwich}) = \{ \text{sandwich} : \text{sandwich} \cdot \text{sandwich}^R \in \text{Sandwich} \},$$

where sandwich^R denote the reversal of the string *sandwich*.

* (b) Let *Fish* be an automatic language.

$$\text{Chop}(\text{Fish}) = \left\{ \text{body} : \begin{array}{l} \text{head} \cdot \text{body} \cdot \text{tail} \in \text{Fish} \text{ for some } \text{head} \text{ and } \text{tail}, \text{ and} \\ \text{all three } \text{head}, \text{ body, and } \text{tail} \text{ have the same length} \end{array} \right\},$$

★ (c) Let *SushiRoll* be an automatic language.

$$\text{Cut}(\text{SushiRoll}) = \{ \text{sushi} : \text{sushi}^n \in \text{SushiRoll} \text{ for some } n \geq 0 \},$$

where sushi^n denote the concatenation of the string *sushi* with itself n times.

2. **Prefix codes.** *Huffman code* is an efficient lossless encoding method that achieves optimal (symbol-to-symbol) compression rate when the input probability distribution is known. The *prefix-free* property, that no codeword inside the constructed encoding is the prefix of another codeword, makes Huffman code uniquely decodable and very efficient in practice. Let L be an automatic language. Construct NFAs that recognize the following languages.

(a) $\{w \in \Sigma^* : \text{no proper prefix of } w \text{ is in } L\}$

(b) $\{w \in \Sigma^* : \text{no proper suffix of } w \text{ is in } L\}$

* (c) $\{w \in \Sigma^* : \text{no proper } \textit{substring} \text{ of } w \text{ is in } L\}$